

A triangulated category \mathcal{J} is an additive category (k -linear). This comes with

- Shift functors: $[t]: \mathcal{J} \rightarrow \mathcal{J}$ for each $t \in \mathbb{Z}$ (Note $[t_1] \cdot [t_2] = [t_1 + t_2]$).

- Class of distinguished triangles

$$T_1 \xrightarrow{f_1} T_2 \xrightarrow{f_2} T_3 \xrightarrow{f_3} T_1[1].$$

↑ better to have equality and not \simeq . Can always pass to an equiv. category to make this equality.

The reader is referred to Gelfand-Manin for all axioms of triangulated categories. The most important axiom is:

For any $f: T_1 \rightarrow T_2$, \exists a distinguished triangle $T_1 \xrightarrow{f} T_2 \rightarrow T_3 \rightarrow T_1[1]$

$T_3 =: \text{Cone}(f)$.

The isomorphism class of $\text{Cone}(f)$ is well defined but is not functorial.

Some remarks about

$$\begin{array}{ccccccc} T_1' & \rightarrow & T_2' & \rightarrow & T_3' & \rightarrow & T_1'[1] \\ \downarrow & & \downarrow & & \downarrow & & \\ T_1 & \xrightarrow{f} & T_2 & \rightarrow & T_3 & \rightarrow & T_1[1] \end{array}$$

Some notation:

- $D(\mathcal{C}) :=$ the derived category of the abelian category \mathcal{C} .

- $D^b(\mathcal{C}) =$ Bounded derived category.

- $\text{Ext}^i(T, T') = \text{Hom}^i(T, T') = \text{Hom}(T, T'[i])$

- $\text{Hom}^*(T, T') = \bigoplus_{i \in \mathbb{Z}} \text{Hom}^i(T, T')[-i] \in \text{grmod}(k) = D(k)$.

- $\forall v \in \text{mod}(k); T \in \mathcal{J}$ then $V \otimes_k T = T^{\oplus \dim V}$

This is not a good definition (non-functorial)

So you define it by the functor it represents.

$$\left. \begin{array}{l} \text{Hom}(V \otimes T, T') = \text{Hom}(V^*, \text{Hom}(T, T')) \\ \text{Hom}(T', V \otimes T) = V \otimes \text{Hom}(T', T) \end{array} \right\} \textcircled{*}$$

- $V^* \in \text{grmod}(k); V^* = \bigoplus_{i \in \mathbb{Z}} V_i[-i]$ (Replace Hom with Hom^*)

$V^* \otimes T = \bigoplus_{i \in \mathbb{Z}} V_i \otimes T[-i]$ and V with V^* in the above conditions. (*)

The most important notion for these lectures is that of:

Semiorthogonal Decomposition

Let \mathcal{J} be a triangulated category, $\mathcal{A}, \mathcal{B} \subset \mathcal{J}$ be strictly full triangulated subcategories

is closed under shifts and cones.

Def: $\mathcal{J} = \langle \mathcal{A}, \mathcal{B} \rangle$ is a semiorthogonal decomposition if:

Can also take Hom^*

1) $\text{Hom}(\mathcal{B}, \mathcal{A}) = 0$ (ie $\forall A \in \mathcal{A}, B \in \mathcal{B}$ we have $\text{Hom}(B, A) = 0$)

2) For any $T \in \mathcal{J}$ there is a distinguished triangle $T_{\mathcal{B}} \rightarrow T \rightarrow T_{\mathcal{A}} \rightarrow T_{\mathcal{B}}[1]$.

with $T_{\mathcal{A}} \in \mathcal{A}$ and $T_{\mathcal{B}} \in \mathcal{B}$.

Properties:

1) Functoriality: For any $f: T \rightarrow T'$ there exist unique $f_{\mathcal{A}}$ and $f_{\mathcal{B}}$ such that the following

diagram commutes:

$$\begin{array}{ccccccc}
 & & & \text{in proof} & & & \\
 & & & \dots \rightarrow \text{map below is just composition with this arrow} & & & \\
 T_{\mathcal{B}} & \rightarrow & T & \rightarrow & T_{\mathcal{A}} & \rightarrow & T_{\mathcal{B}}[1] \\
 f_{\mathcal{B}} \downarrow & & \downarrow f & & \downarrow f_{\mathcal{A}} & & \downarrow f_{\mathcal{B}}[1] \\
 T'_{\mathcal{B}} & \rightarrow & T' & \rightarrow & T'_{\mathcal{A}} & \rightarrow & T'_{\mathcal{B}}[1] \\
 & & & & \text{comments in proof} & &
 \end{array}$$

Proof: Apply $\text{Hom}(-, T'_{\mathcal{A}})$. Get long exact sequence

$$\dots \rightarrow \text{Hom}(T_{\mathcal{B}}[1], T'_{\mathcal{A}}) \rightarrow \text{Hom}(T_{\mathcal{A}}, T'_{\mathcal{A}}) \xrightarrow{\cong} \text{Hom}(T, T'_{\mathcal{A}}) \rightarrow \text{Hom}(T_{\mathcal{B}}, T'_{\mathcal{A}}) \rightarrow \dots$$

$\begin{matrix} \text{"} & & \text{"} \\ 0 & & 0 \end{matrix}$

$\Rightarrow \exists! f_{\mathcal{A}}$ s.t. the middle square commutes.

The exercise shows that the

Exercise: $\exists! f_{\mathcal{B}}$ s.t. the left square commutes. \square

left square commutes.

Corollary:

$$\left. \begin{array}{l} T \mapsto T_{\mathcal{A}} \\ f \mapsto f_{\mathcal{A}} \end{array} \right\} \text{ is a functor and so is } \left. \begin{array}{l} T \mapsto T_{\mathcal{B}} \\ f \mapsto f_{\mathcal{B}} \end{array} \right\}$$

These are both functors from $\mathcal{J} \rightarrow \mathcal{A}$ and $\mathcal{J} \rightarrow \mathcal{B}$.

Let the adjoint functors be:

$$\mathcal{A} \xrightarrow{\alpha} \mathcal{J} \text{ and } \mathcal{B} \xrightarrow{\beta} \mathcal{J}$$

Then $T \mapsto T_{\mathcal{A}}$ is left adjoint to α (denoted α^*) and $T \mapsto T_{\mathcal{B}}$ is right adjoint to β (denoted $\beta^!$)

$$\text{Hom}(T, \alpha A) \ni f \quad \xrightarrow{\sim} \quad f_{\alpha} \in \text{Hom}_{\mathcal{D}}(T_{\alpha}, \alpha A) = \text{Hom}_{\mathcal{A}}(T_{\alpha}, A)$$

this association gives
an isomorphism as shown

$$\begin{array}{ccccccc} T_{\mathcal{B}} & \rightarrow & T & \rightarrow & T_{\alpha} & \rightarrow & T_{\mathcal{B}}[1] \\ \downarrow & & \downarrow f & & \downarrow f_{\alpha} & & \downarrow \\ 0 & \rightarrow & \alpha A & \xrightarrow{\text{id}} & \alpha A & \rightarrow & 0 \end{array} \quad \begin{array}{l} T_{\alpha} = \alpha^* T \\ f_{\alpha} = \alpha^* f \end{array}$$

Exercise: Prove the above \cong for $T_{\mathcal{B}}$.

Definition: A full subcategory $\mathcal{A} \xrightarrow{\alpha} \mathcal{J}$ is left admissible if $\exists \alpha^*: \mathcal{J} \rightarrow \mathcal{A}$.
right admissible if $\exists \alpha^!: \mathcal{J} \rightarrow \mathcal{A}$.

If $\mathcal{J} = \langle \mathcal{A}, \mathcal{B} \rangle$ then:

$$\text{Exercise: } \begin{cases} \mathcal{B} = {}^{\perp} \mathcal{A} := \{ T \in \mathcal{J} \mid \text{Hom}(T, \mathcal{A}) = 0 \} = \text{Ker } \alpha^* \\ \mathcal{A} = \mathcal{B}^{\perp} := \{ T \in \mathcal{J} \mid \text{Hom}(\mathcal{B}, T) = 0 \} = \text{Ker } \alpha^! \end{cases}$$

Proposition: If $\mathcal{A} \xrightarrow{\alpha} \mathcal{J}$ is left admissible $\Rightarrow \mathcal{J} = \langle \mathcal{A}, {}^{\perp} \mathcal{A} \rangle$
If $\mathcal{B} \hookrightarrow \mathcal{J}$ is right admissible $\Rightarrow \mathcal{J} = \langle \mathcal{B}^{\perp}, \mathcal{B} \rangle$

Proof: $\alpha^*: \mathcal{J} \rightarrow \mathcal{A} \quad T \in \mathcal{J}, \alpha^* T \in \mathcal{A}, \alpha \alpha^* T \in \mathcal{J}$.

aim to show this \rightarrow

$$\begin{array}{c} T' \rightarrow T \xrightarrow{(\ast)} \alpha \alpha^* T \rightarrow T'[1] \\ \uparrow \mathcal{A} \quad \uparrow \mathcal{A} \end{array}$$

$$\begin{array}{ccccccc} \alpha^* T' & \rightarrow & \alpha^* T & \rightarrow & \alpha^* \alpha \alpha^* T & \rightarrow & \alpha^* T'[1] \\ \parallel & & & \searrow & \parallel & & \\ 0 & & & \xrightarrow{\text{id}} & \alpha^* T & & \end{array}$$

from axioms of triangulated categories.

because (\ast) came from adjunction

Definition: Suppose that $\mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_n \in \mathcal{J}$ (strictly full triangulated subcategories)

We say $\mathcal{J} = \langle \mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_n \rangle$ is a semi-orthogonal decomposition if:

1) $\text{Hom}(\mathcal{A}_i, \mathcal{A}_j) = 0$ for $i > j$.

2) For all $T \in \mathcal{J}$ there exists $0 = T_n \rightarrow T_{n-1} \rightarrow \dots \rightarrow T_1 \rightarrow T_0 = T$, such that

$\text{Cone}(T_i \rightarrow T_{i-1}) \in \mathcal{A}_i$

Remark: Set $n=2$. Then

$0 = T_2 \rightarrow T_1 \rightarrow T_0 = T$

you just take the part of the data to get the required triangle

$T_1 \rightarrow T \rightarrow \text{Cone}(T_1 \rightarrow T) \rightarrow T_1[1]$
 $\uparrow \qquad \qquad \qquad \uparrow$
 $\mathcal{A}_2 \qquad \qquad \qquad \mathcal{A}_1$

$\text{Cone}(T_2 \rightarrow T_1) \in \mathcal{A}_2$

$\text{Cone}(T_1 \rightarrow T) \in \mathcal{A}_1$

So this gives the triangle for the $n=2$ definition.

Exercise:

1) Prove functoriality of $T \mapsto \text{Cone}(T_i \rightarrow T_j) \quad \forall i > j$.

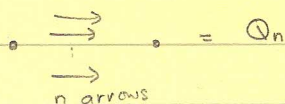
2) $\mathcal{J} = \langle \mathcal{A}_1, \dots, \mathcal{A}_n \rangle \Rightarrow \mathcal{J} = \langle \mathcal{A}_{\leq i}, \mathcal{A}_{\geq i+1} \rangle$ where $\mathcal{A}_{\leq i} = \langle \mathcal{A}_1, \dots, \mathcal{A}_i \rangle$ and $\mathcal{A}_{\geq i+1} = \langle \mathcal{A}_{i+1}, \dots, \mathcal{A}_n \rangle$.

3) $\mathcal{A}_i = \langle \mathcal{A}_1, \dots, \mathcal{A}_{i-1} \rangle^\perp \cap {}^\perp \langle \mathcal{A}_{i+1}, \dots, \mathcal{A}_n \rangle$

4) If you have a semi-orthog. collection of triangulated subcategories, $\mathcal{A}_1, \dots, \mathcal{A}_n \subset \mathcal{J}$ and $\mathcal{A}_1, \dots, \mathcal{A}_i$ are left admissible and $\mathcal{A}_{i+1}, \dots, \mathcal{A}_n$ are right admissible.

Then $\mathcal{J} = \langle \mathcal{A}_1, \dots, \mathcal{A}_i, \mathcal{B}, \mathcal{A}_{i+1}, \dots, \mathcal{A}_n \rangle$
 $\mathcal{A}_{\leq i}^\perp \cap {}^\perp \mathcal{A}_{\geq i+1}$

Example: Take the Kronecker quiver with 2 verts and n arrows:



Then $D(\text{Rep } Q_n) = \langle S_1, S_2 \rangle$. (more on this tomorrow)

(end.)

Lecture 2.

How to check that a subcategory is admissible?

Saturatedness.

To define this we need the notion of covariant cohomological functors.

A covariant cohomological functor $\mathcal{J} \rightarrow \text{mod-}k$ is a functor $H: \mathcal{J} \rightarrow \text{mod-}k$ such that for any distinguished triangle $T_1 \rightarrow T_2 \rightarrow T_3 \rightarrow T_1[1]$ in \mathcal{J} , we have that

$$\dots \rightarrow H(T_1) \rightarrow H(T_2) \rightarrow H(T_3) \rightarrow H(T_1[1]) \rightarrow \dots$$

is a long exact sequence.

Example: For each $T \in \mathcal{J}$ we have representable functors:

$$h_T: \mathcal{J} \rightarrow \text{mod-}k; \quad h_T(-) = \text{Hom}(T, -)$$

$$h^T: \mathcal{J}^{\text{op}} \rightarrow \text{mod-}k; \quad h^T(-) = \text{Hom}(-, T)$$

Definition: A category \mathcal{J} (of finite type) is right (left) saturated if each contravariant (covariant) cohomology functor is representable.

Theorem: • X a smooth proj. variety $\Rightarrow D^b(X)$ is saturated.

- A is a finite dimensional algebra of finite global dimension then $D^b(\text{mod-}A)$ is saturated.
- If \mathcal{J} is saturated, and $\mathcal{A} \subseteq \mathcal{J}$ is left (right) admissible. Then \mathcal{A} is saturated.
- $\mathcal{J} = \langle \mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_n \rangle$ with all \mathcal{A}_i are saturated. Then \mathcal{J} is saturated.
- If \mathcal{J} has a strong generator then \mathcal{J} is saturated.

\nearrow
This implies all the above.

Exercise: Prove by hand (not using the above Theorem) that $D^b(\text{mod-}k)$ is saturated. \swarrow vect. spaces

Proposition: Let $\mathcal{A} \xrightarrow{\alpha} \mathcal{J}$ be a full triangulated subcategory.

If \mathcal{A} is left (resp. right) saturated $\Rightarrow \mathcal{A}$ is left (right) admissible.

Proof: $T \in \mathcal{J}$, $T_B \rightarrow T \rightarrow T_A \rightarrow T_B[1]$ where $\mathcal{B} = {}^{\perp}\mathcal{A} \ni T'$
 $\text{Hom}(T', \alpha(-)) = 0$
 $h_{T'} \circ \alpha = 0$

$$h_T \circ \alpha : \mathcal{A} \rightarrow \text{mod-}k$$

" $\text{Hom}(T, \alpha(-)) \cong h_A$ for some $A \in \mathcal{A}$ (by saturatedness)

$$\text{Hom}(T, \alpha A) \cong h_A(A) = \text{Hom}(A, A) \cong \mathbb{1}_A$$

$$T' \rightarrow T \rightarrow \alpha A \rightarrow T'[1]$$

$$h_{T'} \circ \alpha \rightarrow h_T \circ \alpha \xrightarrow{\cong} h_{\alpha A} \circ \alpha \Rightarrow h_{T'} \circ \alpha = 0$$

"

" $\text{Hom}(\alpha A, \alpha(-))$

" $\text{Hom}(A, -) = h_A(-)$

□

remark:
 can also prove
 by showing
 $\alpha^*(T) = A$

We now consider the question:

How can one construct a fully faithful functor $D^b(\text{mod-}k) \rightarrow \mathcal{J}$?

" $\text{gr-mod-}k$

$$\bigoplus_{i \in \mathbb{Z}} k^n[i]$$

ie \mathbb{Z}

$$k \longmapsto E$$

In other words, let me take

$$E \in \mathcal{J} \rightsquigarrow \varphi_E : D^b(\text{mod-}k) \rightarrow \mathcal{J}$$

$$V^\bullet \longmapsto V^\bullet \otimes_k E$$

To check
 whether a
 functor is

$$\text{Hom}(\varphi_E(V^\bullet), F) = \text{Hom}(V^\bullet \otimes_k E, F) = \text{Hom}(V^\bullet, \text{Hom}^\bullet(E, F))$$

$$\varphi_E^!(F) = \text{Hom}^\bullet(E, F)$$

$$\varphi_E^*(F) = \text{Hom}^\bullet(F, E)^\vee$$

$$\varphi_E^! \circ \varphi_E(V^\bullet) = \text{Hom}^\bullet(E, V^\bullet \otimes_k E) = \text{Hom}^\bullet(E, E) \otimes V^\bullet$$

φ_E is fully faithful $\Leftrightarrow \text{Hom}^\bullet(E, E) = k$ ie $\text{Hom}(E, E) = k$, $\text{Ext}^i(E, E) = 0$ for $i \neq 0$.

Objects with this useful property have a name (next page):

fully faithful, it
 is useful to
 write down an
 adjoint functor.

Definition: We say an object $E \in \mathcal{T}$ is exceptional if $\text{Hom}^*(E, E) = k$.

Proposition: If $E \in \mathcal{T}$ is exceptional then

$$\bullet \mathcal{T} = \langle \mathcal{U}_E(D^b(\text{mod-}k)), {}^\perp E \rangle$$

$$\bullet \mathcal{T} = \langle E^\perp, \mathcal{U}_E(D^b(\text{mod-}k)) \rangle$$

as an abuse of notation we sometimes write

$$\mathcal{T} = \langle E, {}^\perp E \rangle \text{ and } \mathcal{T} = \langle E^\perp, E \rangle.$$

Definition: An exceptional collection in \mathcal{T} is a collection E_1, E_2, \dots, E_n of exceptional objects which is semi-orthogonal i.e. $\text{Hom}^*(E_i, E_j) = 0$ for $i > j$.

Proposition: For any exceptional collection E_1, \dots, E_n , there is a semi-orthogonal decomposition:

$$\mathcal{T} = \langle E_1, \dots, E_i, {}^\perp \langle E_1, \dots, E_i \rangle \cap \langle E_{i+1}, \dots, E_n \rangle^\perp, E_{i+1}, \dots, E_n \rangle.$$

Definition: An exceptional collection E_1, \dots, E_n is full if it generates \mathcal{T}

$$\text{i.e. } \mathcal{T} = \langle E_1, E_2, \dots, E_n \rangle.$$

Example: $D^b(\mathbb{P}^1)$. We claim that $(\mathcal{O}(-1), \mathcal{O})$ is a full exceptional collection.

Proof: Step 1.

$$\text{Hom}^*(\mathcal{O}(k), \mathcal{O}(l)) = H^*(\mathbb{P}^1, R\mathcal{H}om(\mathcal{O}(k), \mathcal{O}(l)))$$

$$= H^*(\mathbb{P}^1, \mathcal{O}(k)^\vee \otimes^{\mathbb{L}} \mathcal{O}(l))$$

$$= H^*(\mathbb{P}^1, \mathcal{O}(-k) \otimes^{\mathbb{L}} \mathcal{O}(l))$$

$$= H^*(\mathbb{P}^1, \mathcal{O}(l-k))$$

You can try the same calculation

✓ on an arbitrary variety

(*)

In particular any line bundle on a Fano variety is exceptional

For $l = k$, $\text{Hom}(\mathcal{O}(k), \mathcal{O}(k)) = k \Rightarrow$ any line bundle on \mathbb{P}^1 is exceptional

(The same argument works on any connected X s.t. $h^{0,q}(X) = 0$ for $q \geq 1$).

(Cont. \rightarrow)

Have to assume $\text{char}(k) = 0$ to use Kodaira vanishing.

Step 2:
 $\text{Hom}^*(\mathcal{O}_X, \mathcal{O}_X(-iH)) = H^*(\mathbb{P}^1, \mathcal{O}(-i)) = 0$.

If X is a Fano variety then
 $-K_X = r \cdot H$ for $r \in \mathbb{Z}_{>0}$ with H ample.

Then
 $H^*(X, \mathcal{O}_X(-iH)) = 0$ for $1 \leq i \leq r-1$.
 $\Rightarrow (\mathcal{O}_X((1-r)H), \dots, \mathcal{O}_X(-H), \mathcal{O}_X)$ is an exceptional collection in $D^b(X)$.

Step 3: Checking fullness (this is the most involved step typically)

It is enough to check that $\langle \mathcal{O}(-1), \mathcal{O} \rangle^+ = 0$.

We give 3 proofs of this fact.

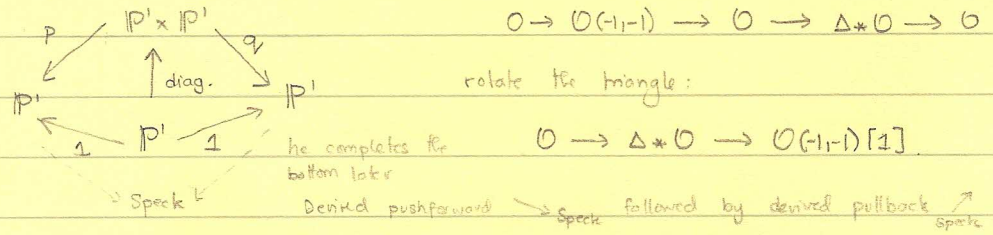
(a) (Very specific to \mathbb{P}^1)

It is enough to check that there are no indecomposables in $\langle \mathcal{O}(-1), \mathcal{O} \rangle^+ = 0$.

$\mathcal{O}(k), \mathcal{O}_X$.

Exercise: Check that these are not in $\langle \mathcal{O}(-1), \mathcal{O} \rangle^+$.

(b) Method of Resolution of the diagonal.



$K \in D^b(\mathbb{P}^1 \times \mathbb{P}^1)$

"Fourier-Mukai functor"

$$\Phi_K : D^b(\mathbb{P}^1) \rightarrow D^b(\mathbb{P}^1)$$

$$F \longmapsto Rq_*(K \otimes^L Lp^* F)$$

Apply this to get:

$$\Phi_0(F) \rightarrow \Phi_{\Delta_* \mathcal{O}}(F) \rightarrow \Phi_{\mathcal{O}(-1,-1)[1]}(F) \quad \text{a distinguished triangle in } D^b(\mathbb{P}^1)$$

(later we conclude that

$$\Phi_0(F) \in \text{Im } \Phi_0 \quad \Phi_{\Delta_* \mathcal{O}}(F) = F \text{ and last term } \in \text{Im } \Phi_{\mathcal{O}(-1,-1)}$$

$$\Phi_0(F) = Rq_*(O \otimes^L Lp^*F) = Rq_*Lp^*F = H^*(IP', F) \otimes O_{P'}.$$

$$\Phi_{O(-1)[1]}(F) = H^*(IP', F(-1)) \otimes O_{P'}(-1)[1].$$

$$\begin{aligned}\Phi_{\Delta^*O}(F) &= Rq_*(R\Delta_*(O \otimes^L Lp^*F)) = Rq_*R\Delta_*(O \otimes^L L\Delta^*Lp^*F) \\ &= R(q \circ \Delta)_*L(p \circ \Delta)^*F = R1_*L1^*F = F.\end{aligned}$$

This completes the proof. (third method next lecture)

(end)

Lecture 3.

We have already discussed two proofs of fullness.

(a) Classification of indecomposables.

(b) Resolution of the diagonal. This technique extends to \mathbb{P}^n and to $\text{Gr}(k, n)$.

(c) In the third approach we will use the Euler sequence:

$$\mathcal{O}(-2) \xrightarrow{\begin{pmatrix} x \\ -y \end{pmatrix}} \mathcal{O}(-1) \oplus \mathcal{O}(-1) \xrightarrow{(x, y)} \mathcal{O}$$

This shows that $\mathcal{O}(-2) \in \langle \mathcal{O}(-1), \mathcal{O} \rangle$. Then look at

$$\mathcal{O}(-3) \rightarrow \mathcal{O}(-2) \oplus \mathcal{O}(-2) \rightarrow \mathcal{O}(-1)$$

Therefore $\mathcal{O}(-3) \in \langle \mathcal{O}(-2), \mathcal{O}(-1) \rangle \subseteq \langle \mathcal{O}(-1), \mathcal{O} \rangle$

By induction we see that $\mathcal{O}(-t) \in \langle \mathcal{O}(-1), \mathcal{O} \rangle$ for all $t \geq 0$. (*)

So it is enough to check that $\langle \mathcal{O}(-1), \mathcal{O} \rangle^\perp = 0$

\cap by *

$$\mathcal{O}(-t)^\perp$$

So it is enough to show that $\cap \mathcal{O}(-t)^\perp = 0$.

Now

$$\text{Hom}^i(\mathcal{O}(-t), F) = H^0(\mathcal{H}^i(F)(t)) \quad \text{for } t \gg 0. \quad \text{then } \mathcal{H}^i(F) = \frac{\text{Ker}(F^i \rightarrow F^{i+1})}{\text{Im}(F^{i-1} \rightarrow F^i)}$$

If $\text{LHS} = 0 \Rightarrow \mathcal{H}^i(F) = 0 \Rightarrow F = 0$, so we are done

Another version of (c):

$$(c)' \quad 0 \rightarrow \mathcal{O}(-1) \rightarrow \mathcal{O} \rightarrow \mathcal{O}_p \rightarrow 0.$$

$$\mathcal{O}_p \in \langle \mathcal{O}(-1), \mathcal{O} \rangle \Rightarrow \langle \mathcal{O}(-1), \mathcal{O} \rangle^\perp \subseteq \mathcal{O}_p^\perp.$$

$$\mathcal{O}_p^\perp = \{F \mid \text{supp } \mathcal{H}^i(F) \not\ni p\}$$

$\therefore \cap \mathcal{O}_p^\perp = \{F \mid \text{supp } \mathcal{H}^i(F) = \emptyset\} = 0$ So we are done.

Examples of full exceptional collections.

1) $D^b(\mathbb{P}^n) = \langle \mathcal{O}(-n), \mathcal{O}(-n+1), \dots, \mathcal{O}(-1), \mathcal{O} \rangle$, the Beilinson exceptional collection.

Exercise: Prove this using methods (b), (c) and (c')

2) $D^b(\mathbb{Q}^n) = \langle \mathcal{A}, \mathcal{O}(1-n), \dots, \mathcal{O}(-1), \mathcal{O} \rangle$
 \mathbb{P}^{n+1} quadric

In the case when k is algebraically closed of char $k = 0$,

then $\mathcal{A} = \begin{cases} \langle S \rangle & \text{if } n \text{ is odd (} S \text{ is called a spinor bundle)} \\ \langle S^+, S^- \rangle & \text{if } n \text{ is even (} S^+, S^- \text{ are spinor bundles)} \end{cases}$

For instance:

$n=1, S = \mathcal{O}(-1)$

$n=2, S^+, S^- = \mathcal{O}(-1, -2), \mathcal{O}(-2, -1)$

$n=4, S^\pm$ - twists of tautological bundles.

3) $D^b(\text{Gr}(k, n))$

Let $\mathcal{U} \subseteq \mathcal{O}^{\oplus n}$ be the tautological bundle. $r(\mathcal{U}) = k$.

$\mathcal{U} \rightarrow \mathcal{G}$ - $\text{GL}(k)$ -bundle.

For any sequence of integers $(\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_k)$ we have $\sum^\lambda(\mathcal{U})$ (Schur functor)

eg. $\lambda = (a, 0, \dots, 0) \Rightarrow \sum^\lambda(\mathcal{U}) = S^a \mathcal{U}$; $\sum^{(0, \dots, 0)} \mathcal{U} = \mathcal{O}$

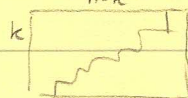
$\lambda = (\underbrace{1, \dots, 1}_a, \underbrace{0, \dots, 0}_{k-a}) \Rightarrow \sum^\lambda \mathcal{U} = \Lambda^a \mathcal{U}$; $\sum^{(1, \dots, 1)} \mathcal{U} = \Lambda^k \mathcal{U} = \mathcal{O}(-1)$

We have

$D^b(\text{Gr}(k, n)) = \langle \sum^\lambda(\mathcal{U}) \rangle_{n-k \geq \lambda_1 \geq \dots \geq \lambda_k \geq 0}$ ← this exceptional collection is indexed by

the order is opposite to "C"

Young diagrams



Let k be alg. closed of char $= 0$.

Conjecture: For any semisimple G and any parabolic PCC, there is a full exceptional collection in $D^b(G/P)$.

Known cases:

- 1) G is of type A
- 2) types B, C, D it is known for some G/P . Also known for flag varieties
- 3) G_2 - known
- 4) E_6, F_4 known for some G/P .
 E_7, E_8 nothing is known.

Exercise: If E_1, E_2, \dots, E_n is a full exceptional collection in an arbitrary triangulated category \mathcal{T} then $\{[E_1], [E_2], \dots, [E_n]\}$ is a \mathbb{Z} -basis in $K_0(\mathcal{T})$
(So we get an obstruction to the existence of a full exceptional collection).

Exercise: Show that any full exceptional collection in $D^b(\mathbb{P}^1)$ is:

$$(O(k)[s], O(k+1)[t]) \text{ for some } k, s, t \in \mathbb{Z}.$$

Mutation of Exceptional Collections.

Suppose $\mathcal{T} = \langle \mathcal{A}, \mathcal{B} \rangle$ then we can write down also $\mathcal{T} = \langle \mathcal{C}, \mathcal{E} \rangle$

where $\mathcal{A} = \mathcal{B}^\perp$
you should think of Hom as a type of pairing

$$\mathcal{E} = {}^\perp \mathcal{B}$$

Let us write down an explicit equivalence.

$$\alpha, \beta, \gamma$$

$$\alpha^*, \beta^*, \beta', \gamma'$$

Proposition: $\mathcal{A} \cong \mathcal{E}$. More precisely $\gamma' \alpha$ and $\alpha^* \gamma$ are mutually inverse

Proof: Start with an arbitrary object $T \in \mathcal{T}$. Then the composition triangle for T looks like:

$$\gamma \gamma' T \rightarrow T \rightarrow \beta \beta^* T \quad \text{where the morphisms are canonical}$$

Set $T = \alpha A$

$$\gamma \gamma' \alpha A \rightarrow \alpha A \rightarrow \beta \beta^* \alpha A$$

$$\alpha^* \beta = 0 \text{ by semi-orthogonality}$$

$$\alpha^* \gamma \gamma' \alpha A \rightarrow \alpha^* \alpha A \rightarrow \alpha^* \beta \beta^* \alpha A$$

$$\uparrow \quad \quad \quad \uparrow$$

$A \quad \quad \quad 0 \text{ because } \mathcal{B} = \ker \alpha^*$

therefore $\ker \alpha^* = A$ completing the proof. \square

Corollary:

a) $\alpha\alpha^*: \mathcal{J} \rightarrow \mathcal{J}$

1) vanishes on \mathcal{B}

2) induces an equivalence $\mathcal{C} \xrightarrow{\sim} \mathcal{A}$

b) $\gamma\gamma^!: \mathcal{J} \rightarrow \mathcal{J}$

1) vanishes on \mathcal{B}

2) induces an equivalence $\mathcal{A} \rightarrow \mathcal{C}$.

Definition: $L_{\mathcal{B}} := \alpha\alpha^*$ is called the left mutation functor

$R_{\mathcal{B}} := \gamma\gamma^!$ is called the right mutation functor.

Proposition: If $\mathcal{J} = \langle \mathcal{A}_1, \dots, \mathcal{A}_i, \dots, \mathcal{A}_n \rangle$ is a semi-orthogonal decomposition

then for any i , the following are semi-orthogonal decompositions:

$$\mathcal{J} = \langle \mathcal{A}_1, \dots, \mathcal{A}_{i-2}, \mathcal{A}_i, R_{\mathcal{A}_i}(\mathcal{A}_{i-1}), \mathcal{A}_{i+1}, \dots, \mathcal{A}_n \rangle$$

$$\mathcal{J} = \langle \mathcal{A}_1, \dots, \mathcal{A}_{i-1}, L_{\mathcal{A}_i}(\mathcal{A}_{i+1}), \mathcal{A}_i, \mathcal{A}_{i+2}, \dots, \mathcal{A}_n \rangle$$

In general we have the following triangle for right mutation functors

$$R_{\mathcal{B}} \rightarrow \text{id} \rightarrow \beta\beta^* \quad \textcircled{1}$$

and for left mutation functors

$$\beta\beta^! \rightarrow \text{id} \rightarrow L_{\mathcal{B}}$$

So $R_{\mathcal{A}_i}(\mathcal{A}_{i-1}) \subseteq \langle \mathcal{A}_i, \mathcal{A}_i \rangle$

The semiorthogonality of $\mathcal{A}_i, R_{\mathcal{A}_i}(\mathcal{A}_{i-1})$ follows from the corollary

One category lies in $R_{\mathcal{A}_i}(\mathcal{A}_{i-1})$ and the other in \mathcal{A}_i , arguing using $\textcircled{1}$.

Exercise:

1) Compute the mutations $L_{\mathcal{O}(k)}\mathcal{O}(k+1)$ and $R_{\mathcal{O}(k+1)}\mathcal{O}(k)$ on \mathbb{P}^1

2) Do all compositions of at most 3 mutations to $(\mathcal{O}(-2), \mathcal{O}(-1), \mathcal{O})$ on \mathbb{P}^2

Lecture 4.

Last lecture we discussed left and right mutations.

The left mutation is the cone of the morphism $\beta\beta' \rightarrow \text{id}_T$

$$\text{ie } \beta\beta' \rightarrow \text{id}_T \rightarrow L_B$$

and similarly the right mutation is the fiber:

$$R_B \rightarrow \text{id}_T \rightarrow \beta\beta'$$

This gives a braid group action.

To see this, we will check that the braid group relations hold:

Assume WLOG $T = \langle \alpha_1, \alpha_2, \alpha_3 \rangle$

We will show that: $R_1 R_2 R_1 = R_2 R_1 R_2$.

$$\text{Proof: } \langle \alpha_1, \alpha_2, \alpha_3 \rangle \xrightarrow{R_1} \langle \alpha_2, R_{\alpha_1}(\alpha_2), \alpha_3 \rangle \xrightarrow{R_2} \langle \alpha_2, \alpha_3, R_{\alpha_3}(R_{\alpha_2}(\alpha_1)) \rangle$$

$\downarrow R_1$

$$\langle \alpha_3, R_{\alpha_3}(\alpha_2), R_{\alpha_3}(R_{\alpha_2}(\alpha_1)) \rangle$$

①

also

$$\langle \alpha_1, \alpha_2, \alpha_3 \rangle \xrightarrow{R_2} \langle \alpha_1, \alpha_3, R_{\alpha_3}(\alpha_2) \rangle \xrightarrow{R_1} \langle \alpha_3, R_{\alpha_3}(\alpha_1), R_{\alpha_3}(\alpha_2) \rangle$$

$\downarrow R_2$

$$\langle \alpha_3, R_{\alpha_3}(\alpha_2), R_{R_{\alpha_3}(\alpha_2)}(R_{\alpha_3}(\alpha_1)) \rangle$$

②

and these coincide - either by a direct argument

or using the following exercise:

Exercise: $T = \langle \beta_1, \beta_2 \rangle$ Then

$$R_T = R_{\beta_2} R_{\beta_1} \quad \text{and} \quad L_T = L_{\beta_1} \circ L_{\beta_2}$$

Using this exercise:

$$\textcircled{1} = R_{\langle \alpha_2, \alpha_3 \rangle}(\alpha_1) \quad \text{and} \quad \textcircled{2} = R_{\langle \alpha_3, R_{\alpha_3}(\alpha_2) \rangle}(\alpha_1)$$

Serre Functors

These are categorical interpretations of Serre Duality.

Definition: A Serre functor for \mathcal{T} is an autoequivalence $S: \mathcal{T} \rightarrow \mathcal{T}$ such that

$$\text{Hom}^\circ(\mathcal{T}_1, \mathcal{T}_2)^\vee \cong \text{Hom}^\circ(\mathcal{T}_2, S\mathcal{T}_1).$$

Example: If $\mathcal{T} = D^b(X)$ with X smooth, projective then $S = - \otimes \omega_X[\dim X]$ is a Serre functor.

For instance: if $D^b(\text{mod-}k) = D^b(\text{Spec } k) \Rightarrow S = \text{id}$.

Proposition: If a Serre functor exists, then it is unique.

Proof: $S_1, S_2: \mathcal{T} \rightarrow \mathcal{T}$. Then

$$\text{Hom}(S_1 T, S_2 T) = \text{Hom}(T, S_1 T)^\vee = \text{Hom}(T, T)^{\vee\vee} = \text{Hom}(T, T) \ni \text{id}_T. \quad \square$$

Proposition: If \mathcal{T} is saturated then $S_{\mathcal{T}}$ exists.

Proof: $h_{\mathcal{T}}^\vee \cong h^{S\mathcal{T}_1}$

□

Proposition: $\langle \mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_n \rangle$. Then

$$L_{\mathcal{A}_1} L_{\mathcal{A}_2} \dots L_{\mathcal{A}_{n-1}}(\mathcal{A}_n) = S_{\mathcal{T}}(\mathcal{A}_n) \text{ and}$$

$$R_{\mathcal{A}_n} R_{\mathcal{A}_{n-1}} \dots R_{\mathcal{A}_2}(\mathcal{A}_1) = S_{\mathcal{T}}^{-1}(\mathcal{A}_1).$$

Proof: $L_{\mathcal{A}_1} \dots L_{\mathcal{A}_{n-1}}(\mathcal{A}_n) = \langle \mathcal{A}_1, \dots, \mathcal{A}_{n-1} \rangle^\perp$

$$\mathcal{A}_n = {}^\perp \langle \mathcal{A}_1, \dots, \mathcal{A}_{n-1} \rangle$$

$$\mathcal{B} \subseteq \mathcal{T}$$

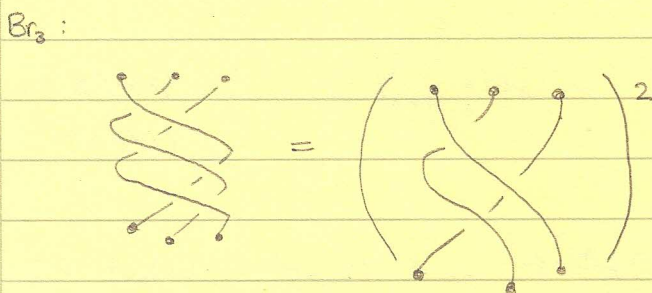
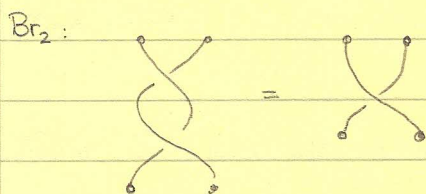
$$S_{\mathcal{T}}({}^\perp \mathcal{B}) = \mathcal{B}^\perp \text{ and}$$

$$S_{\mathcal{T}}^{-1}(\mathcal{B}^\perp) = {}^\perp \mathcal{B}. \quad \square$$

Now it is well known that $Z(\text{Br}_n) \cong \mathbb{Z}$. The generator is a distinguished ^{central} element of the braid group $(L_1, L_2, \dots, L_{n-1})^n$ and it just acts as the Serre functor S_J because:

$$\begin{aligned}
 J &= \langle \alpha_1, \dots, \alpha_n \rangle \\
 &\quad \downarrow L_1 \dots L_{n-1} \\
 &\langle S_J(\alpha_n), \alpha_1, \dots, \alpha_{n-1} \rangle \\
 &\quad \downarrow L_1 \dots L_{n-1} \\
 &\langle S_J(\alpha_{n-1}), S_J(\alpha_n), \alpha_1, \dots, \alpha_{n-2} \rangle \\
 &\quad \downarrow \\
 &\quad \vdots \\
 &\quad \downarrow \\
 &\langle S_J(\alpha_1), S_J(\alpha_2), \dots, S_J(\alpha_n) \rangle
 \end{aligned}$$

Similarly, $(R_{n-1}, \dots, R_2, R_1)^n = S_J^{-1}$



$$\begin{aligned}
 \text{Br}_n \\
 \cup \\
 D_n &= |D_{n-1} \circ R_{n-1} \circ \dots \circ R_1 \\
 D_n^2 &= (|R_{n-1} \dots |R_1)^n
 \end{aligned}$$

Exercise: $\langle \mathcal{A}_1, \dots, \mathcal{A}_n \rangle \xrightarrow{D_n} \langle \mathcal{B}_1, \dots, \mathcal{B}_2 \rangle$ Right Dual Semi-Orthogonal Collection.

$$\mathcal{B}_i = {}^\perp \langle \mathcal{A}_1, \dots, \mathcal{A}_{i-1}, \mathcal{A}_{i+1}, \dots, \mathcal{A}_n \rangle$$

ie. $\text{Hom}(\mathcal{B}_i, \mathcal{A}_j) = 0$ for $i \neq j$.

$$\text{If } \langle E_1, \dots, E_n \rangle \xrightarrow{D_n} \langle F_n, \dots, F_1 \rangle$$

Right dual exceptional collection.

$$\text{Hom}(E_i, E_j) = \begin{cases} 0 & \text{if } i \neq j \\ k & \text{if } i = j \end{cases}$$

This can be used to construct:

Resolution of the Diagonal.

X is smooth projective. Consider $D^b(X)$.

$$D^b(X) = \langle E_1, E_2, \dots, E_n \rangle$$

$$D^b(X) = \langle F_n, F_{n-1}, \dots, F_1 \rangle$$

$$X \times X \xleftarrow{\Delta} X$$

Theorem: There is a chain of maps:

$$0 = D_n \rightarrow D_{n-1} \rightarrow \dots \rightarrow D_1 \rightarrow D_0 = \Delta_* \mathcal{O}_X, \text{ such that } \text{Cone}(D_i \rightarrow D_{i-1}) = E_i \boxtimes F_i^\vee$$

"sum of rank 1 vectors"

$$k \in D^b(X \times X)$$

$$\Phi_{X \leftarrow X}^k = \Phi_k$$

$$\Gamma \in D^b(X)$$

$$0 = \Phi_{D_n}(T) \rightarrow \Phi_{D_{n-1}}(T) \rightarrow \dots \rightarrow \Phi_{D_1}(T) \rightarrow \Phi_{D_0}(T) = T$$

$$\Phi_{D_i}(T) \rightarrow \Phi_{D_{i-1}}(T) \rightarrow \Phi_{E_i \boxtimes F_i^\vee}(T)$$

$$\text{Hom}^i(F_i, T) \otimes E_i$$

Exercise: Compute the RDEC for:

1) $\langle \mathcal{O}(-n), \dots, \mathcal{O} \rangle$ on \mathbb{P}^n

2) $\langle \Sigma^\lambda \mathcal{U} \rangle$ on $\text{Gr}(k, n)$

Lecture 5.

Recall that last time we discussed Resolution of the Diagonal, which stated the following:

If $D^b(X) = \langle E_1, E_2, \dots, E_n \rangle$, then

$$0 = D_n \rightarrow D_{n-1} \rightarrow \dots \rightarrow D_1 \rightarrow D_0 = \Delta^* O_X \in D^b(X \times X)$$

such that $\text{Cone}(D_i \rightarrow D_{i-1}) = E_i \boxtimes F_i^\vee (= \mathcal{R}\text{Hom}(L_{P_i}^* F_i, L_{P_i}^* E_i))$

to prove this, we need the following bit of theory:

Base Change for Semi-Orthogonal Decompositions.

$$\begin{array}{ccc} X \ni x & & D^b(X) = \langle \mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_n \rangle \\ \downarrow P & & \\ X \xrightarrow{f} S & \xrightarrow{\quad} & S \end{array}$$

Definition: We say that $\mathcal{A} \subseteq D^b(X)$ is S-linear if $\mathcal{A} \otimes_{L_P^*} F \subseteq \mathcal{A}$ for any $F \in D^b(S)$.

$\begin{array}{ccc} X_T \xrightarrow{\tilde{P}} X \\ \downarrow P_T & & \downarrow P \\ T \xrightarrow{f} S \end{array}$	<p>Theorem: Assume that $D^b(X) = \langle \mathcal{A}_1, \dots, \mathcal{A}_n \rangle$ - S-linear semi-orthogonal decomposition.</p> <p>Let f be a base change such that T and X are <u>Tor-independent</u> over S.</p> <p>Then $D^b(X_T) = \langle \mathcal{A}_{1T}, \mathcal{A}_{2T}, \dots, \mathcal{A}_{nT} \rangle$ is a T-linear semi-orthog. decomposition s.t</p> <p>$\mathcal{R}\tilde{F}_*(\mathcal{A}_{iT}) \subseteq \mathcal{A}_i; L\tilde{F}^*(\mathcal{A}_i) \subseteq \mathcal{A}_{iT}$.</p>
--	---

means $\text{Tor}_{>0}^{O_{S,S}}(O_{T,T}, O_{X,X}) = 0$

Idea of Proof:

$$\langle L\tilde{F}^*(\mathcal{A}_i) \otimes_{L_{P_T}^*} D^b(T) \rangle \cong \mathcal{A}_{iT}$$

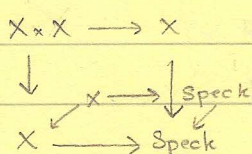
Remark: \mathcal{A}_{iT} does not depend on the choice of embedding $\mathcal{A}_i \subseteq D^b(X)$.

eg. If $\mathcal{A}_i \cong D^b(Y)$ for some Y/S , then $\mathcal{A}_{iT} \subseteq D^b(Y_T) \mid \bigcap^{\cap} D^b(Y)$

(such that Y and T are tor-independent over S .)

Proof of Resolution of the diagonal:


$$D^b(X) = \langle \mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_n \rangle$$



$$D^b(X \times X) = \langle \mathcal{A}_{1X}, \dots, \mathcal{A}_{nX} \rangle$$

$$\Delta^* \mathcal{O}_X$$

$$\text{Cone}(D_i \rightarrow D_{i-1}) = E_i \boxtimes F_i' \quad (\text{before } \mathcal{A}_i \text{ crossed to give } F_i')$$

 bit added later ---

$$D^b(\text{Spec } k) \cong \mathcal{A}_i = \langle E_i \rangle$$

$$D^b(X) \cong \mathcal{A}_{iX} = \{ E_i \boxtimes F \} ; F \in D^b(X)$$

Exercise: Apply the resolution of the diagonal to check that $F_i' \cong F_i^\vee$.

Let $X \xrightarrow{P} S \ni s$ be a flat morphism.

Let $E_1, E_2, \dots, E_n \in D^b(X)$ be such that

$E_{1s}, E_{2s}, \dots, E_{ns} \in D^b(X_s)$ is an exceptional collection for any $s \in S$.

Then:

1) $\varphi_{E_i}: D^b(S) \rightarrow D^b(X)$, $F \mapsto E_i \otimes^L Lp^*(F)$ is fully faithful, $\varphi_{E_i}(D^b(S))$ is admissible and S -linear.

2) $\varphi_{E_1}(D^b(S)), \dots, \varphi_{E_n}(D^b(S))$ is semi-orthogonal.

3) If $E_{1s}, E_{2s}, \dots, E_{ns}$ is full for all $s \in S$ then (\Rightarrow then each X_s is smooth)

$$D^b(X) = \langle \varphi_{E_1}(D^b(S)), \dots, \varphi_{E_n}(D^b(S)) \rangle$$

Proof: (See next page)

Proof:

$$1) \varphi_{E_i}^!(F) = R p_* \mathcal{R} \text{Hom}(E_i, F) = R p_* (E_i^\vee \otimes F)$$

$$\begin{aligned} \varphi_{E_i}^! \circ \varphi_{E_i}(F) &= R p_* \left(E_i^\vee \otimes (E_i \otimes L p^*(F)) \right) \\ &= R p_* \left((E_i^\vee \otimes E_i) \otimes L p^*(F) \right) \\ &= R p_* (E_i^\vee \otimes E_i) \otimes F \quad \text{by projection formula} \end{aligned}$$

$$\begin{aligned} L_{j_s}^* R p_* (E_i^\vee \otimes E_i) &= R p_{s*} L_{i_s}^* (E_i^\vee \otimes E_i) = R p_{s*} (E_{i_s}^\vee \otimes E_{i_s}) \\ &= H^*(X_s, E_{i_s}^\vee \otimes E_{i_s}) = \text{Hom}^*(E_{i_s}, E_{i_s}) = \mathbb{k} \end{aligned}$$

$$\begin{array}{ccc} X & \xleftarrow{i_s} & X_s \\ P \downarrow & & \downarrow P_s \\ S & \xleftarrow{j_s} & S \end{array}$$

$$\mathcal{O}_S \xrightarrow{\cong} R p_* (E_i^\vee \otimes E_i) \xrightarrow{\text{take the cone of the first morphism}} \mathcal{C}$$

$$\mathcal{O}_X = L p^* \mathcal{O}_S \rightarrow E_i^\vee \otimes E_i$$

we have this canonical morphism which induces a canonical morphism $\mathcal{O}_S \rightarrow$ in the sequence above

$\Rightarrow \varphi_{E_i}$ is fully faithful and right admissible.

$$\begin{array}{ccccc} L_{j_s}^* \mathcal{O}_S & \rightarrow & L_{j_s}^* R p_* (E_i^\vee \otimes E_i) & \rightarrow & L_{j_s}^* \mathcal{C} \\ \parallel & \xrightarrow{id} & \parallel & & \parallel \\ \mathbb{k} & & \mathbb{k} & & 0 \end{array}$$

$$\Rightarrow \mathcal{C} = 0.$$

This proves (1).

$$2) \text{Hom}(\varphi_{E_i}(F'), \varphi_{E_j}(F'')) = \text{Hom}(F', \varphi_{E_i}^! \varphi_{E_j}(F''))$$

want this composition = 0

$$\varphi_{E_i}^! \varphi_{E_j}(F) = R p_* (E_i^\vee \otimes E_j) \otimes F$$

$$L_{j_s}^* R p_* (E_i^\vee \otimes E_j) = \text{Hom}^*(E_{i_s}, E_{j_s}) = 0 \quad \text{for } i > j.$$

$$3) D^b(X) = \langle \mathcal{A}, \varphi_{E_1}(D^b(S)), \dots, \varphi_{E_n}(D^b(S)) \rangle$$

$$F \in \mathcal{A} = \langle \dots \rangle^\perp = \ker \varphi_{E_1}^! \cap \dots \cap \ker \varphi_{E_n}^!$$

$$0 = \varphi_{E_i}^!(F) = R p_* (E_i^\vee \otimes F) \Rightarrow 0 = L_{j_s}^* R p_* (E_i^\vee \otimes F) = \text{Hom}^*(E_{i_s}, F_s)$$

$$\Rightarrow \forall s \in S, F_s = 0 \Rightarrow F = 0.$$

This establishes (3).

Example:

$$X = \mathbb{P}_S(\mathcal{E}), \quad r(\mathcal{E}) = n+1$$

$$\mathcal{O}(-n), \mathcal{O}(-n+1), \dots, \mathcal{O}(-1), \mathcal{O}.$$

$$\Rightarrow D^b(X) = \langle \mathcal{O}(-n) \otimes D^b(S), \mathcal{O}(-n+1) \otimes D^b(S), \dots, \mathcal{O}(-1) \otimes D^b(S), D^b(S) \rangle$$

$X \xrightarrow{p} S$ is a \mathbb{P}^n -fibration (Severi-Brauer varieties).

\rightsquigarrow Azumaya Algebra (locally a matrix algebra)

In étale topology $U \rightarrow S$

$$X_U \cong U \times \mathbb{P}^n = \mathbb{P}_U(\mathcal{E}_U)$$

\mathcal{E}_U is defined only up to a twist by a line bundle.

$\text{End}(\mathcal{E}_U) = \mathcal{E}_U^\vee \otimes \mathcal{E}_U$ is canonically defined and so gives to a sheaf of algebras on S . $\mathcal{R} \rightarrow$ azumaya algebra.

$$\text{Br}(S) \ni \beta \leftarrow \mathcal{R}$$

$$\beta \in H_{\text{ét}}^2(S, \mathcal{O}_S^\times)$$

$$D^b(S, \beta) = D^b(\text{coh}(S, \mathcal{R}))$$

\uparrow
 β -twisted sheaves

Theorem (Bernardara)

$\mathcal{O}(-i) \in D^b(X, p^* \beta^{-i})$ which is a $p^* \beta^{-i}$ -twisted sheaf then

$$D^b(X) = \langle \mathcal{O}(-n) \otimes D^b(S, \beta^n), \mathcal{O}(-n+1) \otimes D^b(S, \beta^{n-1}), \dots, \mathcal{O}(-1) \otimes D^b(S, \beta), D^b(S) \rangle$$

(end.)

Lecture 6.

Blowup

Let X be a (not necessarily smooth) variety and let $Z \subseteq X$ be a subvariety.

Let $\tilde{X} = \text{Bl}_Z(X)$

We have the blowup diagram:

$$\begin{array}{ccc} \tilde{X} & \xleftarrow{i} & E \\ \pi \downarrow & & \downarrow p \\ X & \xleftarrow{\quad} & Z \end{array}$$

If $Z \subseteq X$ is a locally complete intersection then: $E \cong \mathbb{P}_Z(\mathcal{U}_{Z/X})$.

Let $c = \text{codim } Z = r(\mathcal{U}_{Z/X})$.

Theorem: (Orlov).

$$D(\tilde{X}) = \langle \text{Ri}_* (\mathcal{O}_{E/Z}(1-c) \otimes Lp^* D^b(Z)), \dots, \text{Ri}_* (\mathcal{O}_{E/Z}(-1) \otimes Lp^* D^b(Z)), L\pi^* D^b(X) \rangle$$

Example:

$$D(\text{Bl}_{pt} \mathbb{P}^2) = \langle \mathcal{O}_E(-1), \mathcal{O}(-2H), \mathcal{O}(-H), 0 \rangle$$

Exercise:

- 1) Mutate $\mathcal{O}_E(-1)$ to the right.
- 2) Write down the resolution of the diagonal and the associated Beilinson spectral sequence.

Other Birational Transformations.

Conjecture:

- a) If $X \dashrightarrow X'$ is a flip then $D(X') \xrightarrow{\sim} D(X)$.
- b) If $X \dashrightarrow X'$ is a flop then $D(X') \cong D(X)$.

(b) Was proved by Bridgeland in the case $\dim = 3$.

There is a more general conjecture:

Let $X \leftarrow \overset{\sim}{\dashrightarrow} X'$ be a birational isomorphism.

$$\begin{array}{ccc} & & \\ \uparrow & & \uparrow \\ X & \overset{\sim}{\dashrightarrow} & X' \\ & \tilde{X} & \end{array}$$

Conjecture: If $K_{\tilde{X}/X} - K_{\tilde{X}/X'} \geq 0$ then $D^b(X) \hookrightarrow D^b(X')$

Griffiths Component of $D^b(X)$

Let $D(X) = \langle \mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_n \rangle$ be a maximal semi-orthogonal decomposition.

Drop all \mathcal{A}_i which can be embedded (as ^{is actually holds automatically} admissible subcategories) into $D^b(Z)$ with $\dim Z \leq \dim X - 2$ and Z smooth

The Griffiths component is the set of components which are left.

Problems: 1) Why should such a maximal semi-orthog. decomposition exist?

2) Is the Griffiths component (as defined) independent of the choice of maximal semi-orthogonal decomposition.

1) should follow from the following:

Conjecture (Noetherian Property).

If X is smooth and projective then any descending chain $D^b(X) \supset \mathcal{A}_1 \supset \mathcal{A}_2 \supset \dots$ stabilizes.

We also remark that the Jordan-Hölder property is false, by the following example:

$$Q = \left\{ \begin{array}{ccc} \bullet & \xrightarrow{a} & \bullet \\ \bullet & \xrightarrow{b} & \bullet \end{array} \quad \begin{array}{ccc} \bullet & \xrightarrow{c} & \bullet \\ \bullet & \xrightarrow{d} & \bullet \end{array} \quad \left| \quad \begin{array}{l} ca = 0 \\ db = 0 \end{array} \right. \right\} \quad D^b(Q) = \langle S_3, S_2, S_1 \rangle$$

$$E = \left\{ \begin{array}{ccc} \bullet & \xrightarrow{1} & \bullet \\ \bullet & \xrightarrow{0} & \bullet \end{array} \quad \begin{array}{ccc} \bullet & \xrightarrow{0} & \bullet \\ \bullet & \xrightarrow{1} & \bullet \end{array} \right\} \quad D^b(Q) = \langle E, {}^\perp E \rangle$$

Other Examples of Semi-Orthogonal Decompositions.

$$S \xleftarrow{P} X \hookrightarrow \mathbb{P}_S(E)$$

$$L \hookrightarrow S^2 E^V \quad r(E) = n.$$

$$D^b(X) = \langle D^b(S, \mathcal{O}_S), \overset{\text{Clifford Algebra}}{\mathcal{O}(2-n) \otimes L_P^*(D^b(S))}, \dots, \mathcal{O}(-1) \otimes L_P^*(D^b(S)), L_P^*(D^b(S)) \rangle$$

Homological Projective Duality.

$$\mathbb{P}(W) \xrightarrow{\vee_2} \mathbb{P}(S^2 W) \supseteq H.$$

$$X \subseteq \mathbb{P}(W) \times \mathbb{P}(S^2 W^V) \quad \text{deg } X = (2, 1)$$



$$\mathbb{P}(S^2 W^V)$$

So we get:

$$D^b(X) = \langle D^b(\mathbb{P}^2(S^2 W^V), \mathcal{O}_S), \underbrace{D^b(\mathbb{P}(S^2 W^V))}_{(n-2) \text{ times}}, \dots, D^b(\mathbb{P}(S^2 W^V)) \rangle$$

Homological Projective Duality gives us:

$$D^b(X_L) = \langle D^b(L, \mathcal{O}_L), \mathcal{O}(1+2d-n), \dots, \mathcal{O} \rangle; \quad n \geq 2d$$

where

$$\mathbb{P}^{d-1} = L \subseteq \mathbb{P}(S^2 W^V)$$

$$X_L = \mathbb{P}(W) \cap L^\perp \subseteq \mathbb{P}(S^2 W)$$

↑
A complete intersection of d quadrics in $\mathbb{P}(W)$

(end of course).