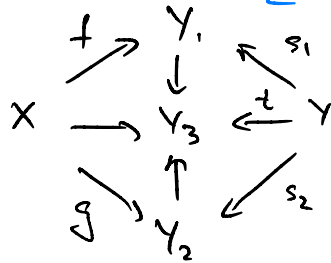




$$\text{Hom}(X, Y) = \left\{ \begin{array}{c} X \xrightarrow{f} Y' \\ \quad \quad \quad \nwarrow \quad \nearrow \\ \quad \quad \quad Y \end{array} \mid s \in S \right\} / \sim$$

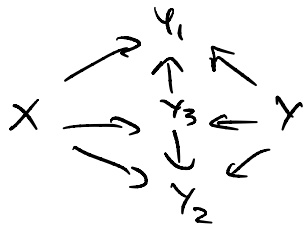
Equivalence relation:



in the previous lecture there is a mistake!  
the vertical arrows were reversed

Need to show that it's an equivalence relation,  
define composition (show that it's well-def).

In order to show transitivity of  $\sim$ , need to show that



gives equivalent morphisms!

Two steps: 1) use  $Y \rightarrow Y_1$

$$\begin{array}{ccc} & & \dots \\ & & \downarrow \\ & & Y_4 \\ Y_2 & \twoheadrightarrow & Y_4 \end{array}$$

2) need to coequalize  $Y \rightarrow Y_3 \rightrightarrows Y_4 \rightarrow Y_5$ .

Application: construct  $\mathcal{A}/\mathcal{B}$ ,  $\mathcal{B}$ -Serre subcategory  
(full, closed under extensions).

Defined  $\mathcal{B}$ -iso's as  $f: A \rightarrow B$  s.t.  $\ker f, \operatorname{coker} f \in \mathcal{B}$ .  
One checks that  $\mathcal{B}$ -iso - both left & right localization  
system.  $\mathcal{A}/\mathcal{B} \neq \mathcal{A}[\mathcal{B}\text{-iso}^{-1}]$  satisfy the same UP.

Q: Why is  $\mathcal{A}[\mathcal{B}\text{-iso}^{-1}]$  abelian? Easiest way - check that  
 $\mathcal{A}/\mathcal{B}$  is abelian (you have explicit descriptions of  
 $\mathcal{B}$ -iso /  $\mathcal{B}$ -mono /  $\mathcal{B}$ -epi).

Finally, we defined  $\mathcal{D}(\mathcal{A}) = \mathcal{C}(\mathcal{A})[\mathcal{Q}is^{-1}]$

category of  $\nearrow$   
complexes

$\nwarrow$  class of  
quasi-isom's

### 1. Complexes in abelian categories

Def  $\mathcal{C}(\mathcal{A})$  consists of objects

$$\dots \rightarrow A^i \xrightarrow{d^{i-1}} A^{i+1} \xrightarrow{d^i} A^{i+2} \xrightarrow{d^{i+1}} \dots$$

$$d^{i+1} \circ d^i = 0 \quad \forall i$$

morphisms

$$\dots \rightarrow A^i \xrightarrow{d} A^{i+1} \rightarrow \dots$$

$$\downarrow f^i \quad \circ \quad \downarrow f^{i+1}$$

$$\dots \rightarrow B^i \rightarrow B^{i+1} \rightarrow \dots$$

composition is obvious.

Exc  $\mathcal{C}(\mathcal{A})$  is abelian if  $\mathcal{A}$  is abelian.

Sketch  $\oplus$  are defined term-wise:  $A^i, B^i \in C(d)$ , then

$$A^i \oplus B^i : \quad \rightarrow A^i \oplus B^i \rightarrow A^{i+1} \oplus B^{i+1} \rightarrow$$

ker & coker are also term-wise!

$$\begin{array}{ccccccc}
 & & \downarrow 0 & & \downarrow 0 & & \\
 \rightarrow & K^i & \rightarrow & K^{i+1} & \rightarrow & & \text{ker } f \\
 & \downarrow & & \downarrow & & & \\
 \rightarrow & A^i & \rightarrow & A^{i+1} & \rightarrow & & \\
 & \downarrow f^i & & \downarrow f^{i+1} & & & \\
 \rightarrow & B^i & \rightarrow & B^{i+1} & \rightarrow & & \\
 & \downarrow & & \downarrow & & & \\
 \rightarrow & C^i & \rightarrow & C^{i+1} & \rightarrow & & \text{Coker } f \\
 & \downarrow & & \downarrow & & & \\
 & 0 & & 0 & & & 
 \end{array}$$

Please, check all the details.

$$A^i \in \mathcal{C}(\mathcal{A}) \rightsquigarrow H^i(A^\bullet) = \ker d^i / \text{Im } d^{i-1}$$

$\ker d^i = Z^i(A^\bullet)$  called cycles

$\text{Im } d^{i-1} = B^i(A^\bullet)$  called boundaries

Get a collection of functors  $H^i: \mathcal{C}(\mathcal{A}) \rightarrow \mathcal{A}$ .

Prop Let  $0 \rightarrow X^\bullet \rightarrow Y^\bullet \rightarrow Z^\bullet \rightarrow 0$  be a short exact sequence in  $\mathcal{C}(\mathcal{A})$ . There exist morphisms  $S: H^i(Z^\bullet) \rightarrow H^{i+1}(Z^\bullet)$  s.t. the following is exact

$$\dots \xrightarrow{S} H^i(X^\bullet) \rightarrow H^i(Y^\bullet) \rightarrow H^i(Z^\bullet) \xrightarrow{S} H^{i+1}(X^\bullet) \rightarrow \dots$$

and given

$$\begin{array}{ccccccc} 0 & \rightarrow & X_1^\bullet & \rightarrow & Y_1^\bullet & \rightarrow & Z_1^\bullet \rightarrow 0 \\ & & \downarrow & & \downarrow & & \downarrow \\ 0 & \rightarrow & X_2^\bullet & \rightarrow & Y_2^\bullet & \rightarrow & Z_2^\bullet \rightarrow 0 \end{array}$$

there is a morphism of LES's  $\Leftrightarrow$

all the squares

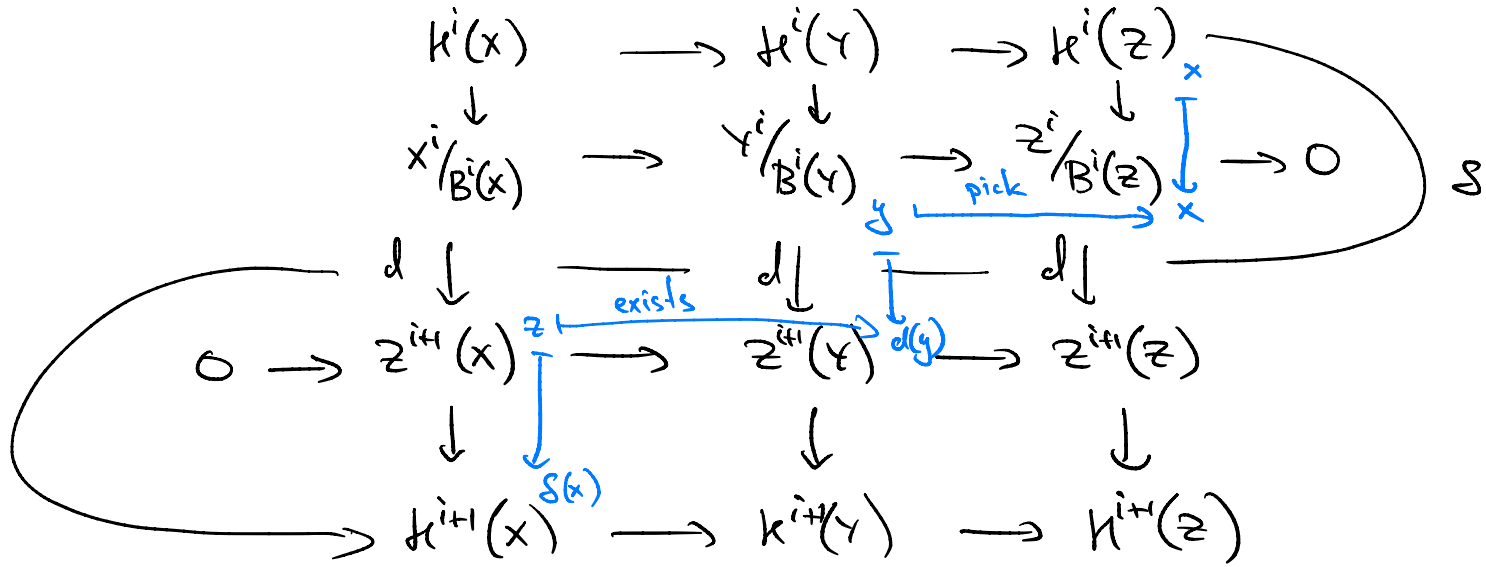
$$\begin{array}{ccc}
 H^i(Z_1) & \xrightarrow{S} & H^{i+1}(X_1) \\
 \downarrow & & \downarrow \\
 H^i(Z_2) & \xrightarrow{S} & H^{i+1}(X_2)
 \end{array}$$

commute.

Pf  $0 \rightarrow X^i \rightarrow Y^i \rightarrow Z^i \rightarrow 0$  is a SES, thus

$$\begin{array}{ccccccc}
 \boxed{0 \rightarrow Z^i(X^i) \rightarrow Z^i(Y^i) \rightarrow Z^i(Z^i)} & & & & & & \\
 \downarrow & \downarrow & \downarrow & & & & \\
 0 \rightarrow X^i \rightarrow Y^i \rightarrow Z^i \rightarrow 0 & & & & & & \\
 \downarrow & \downarrow & \downarrow & & & & \\
 0 \rightarrow X^{i+1} \rightarrow Y^{i+1} \rightarrow Z^{i+1} \rightarrow 0 & & & & & & \\
 \downarrow & \downarrow & \downarrow & & & & \\
 \boxed{X^{i+1}/B^{i+1}(X) \rightarrow Y^{i+1}/B^{i+1}(Y) \rightarrow Z^{i+1}/B^{i+1}(Z) \rightarrow 0} & & & & & & 
 \end{array}$$

Stick these sequences into another 5-lemma type diagram:



How is  $S$  constructed?

Conclusion:

$$\begin{array}{ccccccc}
 0 & \rightarrow & x^i & \rightarrow & y^i & \rightarrow & z^i & \rightarrow & 0 \\
 & & \downarrow & & \downarrow & & \downarrow & & \\
 0 & \rightarrow & x^{i+1} & \rightarrow & y^{i+1} & \rightarrow & z^{i+1} & \rightarrow & 0
 \end{array}$$



- $z \in \mathfrak{t}^i(\mathbb{Z}) \rightsquigarrow$
- 1) pick a representative  $\tilde{z} \in \mathbb{Z}^i$
  - 2) pull it back to  $Y^i$
  - 3) apply  $d_Y^i$
  - 4) pull it back to  $X^{i+1}$ .

□

Recall that long time ago we defined covers of morphisms.

Shift functors  $\Sigma^n: C(\mathcal{A}) \rightarrow C(\mathcal{A})$

$$(X \cdot \Sigma^n)^i = X^{i+n}, \quad d_{X \cdot \Sigma^n} = (-1)^n d_X.$$

Gives an action of  $\mathbb{Z}$  on  $C(\mathcal{A})$ :  $\mathbb{Z} \rightarrow \text{Aut}(C(\mathcal{A}))$ .

Cover  $f: X^i \rightarrow Y^i$  in  $C(\mathcal{A}) \rightsquigarrow C^i(f) \in C(\mathcal{A})$

$$C^i(f) = X^{i+1} \oplus Y^i, \quad d = \begin{pmatrix} -d_X^{i+1} & 0 \\ f^i & d_Y^i \end{pmatrix}$$

Exc Check that  $d^2 = 0$

More imperfectly, there is a SES of complexes:

$$0 \rightarrow Y^\bullet \rightarrow C(f) \rightarrow X[\mathbb{Z}] \rightarrow 0$$

The associated LES of cohomology is

$$\dots \rightarrow H^i(Y^\bullet) \rightarrow H^i(C(f)) \rightarrow H^i(X[\mathbb{Z}]) \xrightarrow{\delta} H^{i+1}(Y) \rightarrow \dots$$

$\downarrow$   
 $H^{i+1}(X)$

Exc Using the explicit description of  $\delta$ , check that for  $0 \rightarrow Y \rightarrow C(f) \rightarrow X[\mathbb{Z}] \rightarrow 0$  the connecting homomorphism is  $\delta = f_*$ , the morphism  $H^i(X) \rightarrow H^i(Y)$ !

Recall that  $f: X \rightarrow Y$  is a quasi-iso if  $H^i(f): H^i(X) \rightarrow H^i(Y)$  - iso  $\forall i \in \mathbb{Z}$ .

We conclude that  $f: X \rightarrow Y$  - quasi-iso  $\Leftrightarrow$   
 $\Leftrightarrow H^i(C(f)) = 0 \quad \forall i.$

$$\begin{array}{ccccccc} & & & H^i(f) & & & \\ \xrightarrow{\sim} & \dots & \rightarrow & H^i(C(f)) & \rightarrow & H^{i+1}(X) & \xrightarrow{\sim} & H^{i+1}(Y) & \rightarrow & H^{i+1}(C(f)) & \rightarrow & \dots \end{array}$$

must be 0!

Think the "analogy" with quotients by Serre subcategories.  
 We want to invert quasi-iso's, looks like it is enough  
 to kill acyclic complexes.

$$\begin{array}{ccccccc} f: X \rightarrow Y & \rightsquigarrow & 0 \rightarrow Y & \rightarrow & C(f) & \rightarrow & X \oplus Y \rightarrow 0 \\ & \uparrow & \text{if qis} & & \uparrow & & \\ & & & \Rightarrow & \text{must go to 0?} & & \end{array}$$

Also remark that  $X^* \in C(\mathcal{A})$  is acyclic (all  $H^i(X^*) = 0$ )  
 iff  $0 \rightarrow X^*$  is a quasi-isomorphism.

Inverting qis's should kill all acyclics.

Problem In  $C(\mathcal{A})$  qis's do not form a localization system (at least, not obvious).

2. Homotopy category

Given  $X^\bullet, Y^\bullet \in C(\mathcal{A})$ , one can produce a complex

$\text{Hom}^\bullet(X^\bullet, Y^\bullet) \in C(\mathcal{A}b)$ :

$\text{Hom}^n(X^\bullet, Y^\bullet) = \prod \text{Hom}(X^i, Y^{i+n}) \leftarrow$  collection of  $X^i \rightarrow Y^{i+n}$ ,  
no relations / commutation with  $d$ 's.

Given  $\varphi \in \text{Hom}^n(X^\bullet, Y^\bullet)$ , put  $d\varphi$  by the rule

$$\begin{array}{ccccccc} \dots & \rightarrow & X^i & \rightarrow & X^{i+1} & \rightarrow & \dots \\ & & \searrow \varphi^i & & \searrow \varphi^{i+1} & & \\ \dots & \rightarrow & Y^{i+n} & \rightarrow & Y^{i+n+1} & \rightarrow & \dots \end{array}$$

$$(d\varphi)^i = d\varphi^i - (-1)^i \varphi^{i+1} d$$

Exc Check that  $\text{Hom}^\bullet(X, Y)$  is indeed a complex.  
Look at its lower terms.

$$Z^0(\text{Hom}^\bullet(X, Y)) = \left\{ \varphi^i: X^i \rightarrow Y^i \mid d\varphi^i - \varphi^{i+1}d = 0 \right\}$$

"  $\text{Hom}_{C(\mathcal{A})}(X, Y)$

$$B^0(\text{Hom}^\bullet(X, Y)) = \left\{ dh^i + h^{i+1}d: X^i \rightarrow Y^i \mid h^i: X^i \rightarrow Y^{i+1} \right\}$$

↗  
morphisms homotopic to zero!

We discussed that there is an equivalence relation on  $\text{Hom}(X, Y)$   $f \sim g \Leftrightarrow f-g$  is homotopic to zero.

Morphisms homotopic to 0 form an ideal:

$$f \sim g \Rightarrow hf \sim hg, fw \sim gw$$

Observation  $f \sim 0 \Rightarrow k^i(f) = 0$  for all  $i \in \mathbb{Z}$ .

Thus,  $f \sim g$ , then  $k^i(f) = k^i(g)$  for all  $i \in \mathbb{Z}$ .

Define  $K(\mathcal{A})$  as the category whose objects are complexes, morphisms - morphisms /  $\sim$ .

$$\text{Hom}_{K(\mathcal{A})}(X, Y) = H^0(\text{Hom}^*(X, Y)).$$

Problem 5 Show that in  $K(\mathcal{A})$  quasi-isomorphisms form a localization system.

$$\begin{array}{ccc}
 X & \xrightarrow{f} & Y \\
 \text{qis } \downarrow s & & \downarrow t \\
 Z & \xrightarrow[\cong]{} & W
 \end{array}
 \quad
 \begin{array}{c}
 \swarrow \\
 tf \sim gs
 \end{array}
 \begin{array}{c}
 \text{homotopic, not} \\
 \text{equal}
 \end{array}$$

Exc Check that  $K(\mathcal{A})[\mathcal{Q}is^{-1}] = \mathcal{D}(\mathcal{A}) = \mathcal{C}(\mathcal{A})[\mathcal{Q}is^{-1}]!$

Hint  $K(\mathcal{A})$  also satisfies some UP.

### 3. Resolutions

Recall that we have a functor  $\mathcal{A} \rightarrow C(\mathcal{A})$ , exact, which sends  $X \mapsto \dots \rightarrow 0 \rightarrow X \rightarrow 0 \rightarrow \dots$   
 $\uparrow$   $n^{\text{th}}$  term

Def A resolution of  $X$  is any complex

$R^\bullet \in C(\mathcal{A})$   $q$  is to  $X$ .

Def A projective resolution of  $X$  - resolution of the form

$\rightarrow P^3 \rightarrow P^2 \rightarrow P^1 \rightarrow P^0 \rightarrow 0 \rightarrow \dots$ ,  $P^i$  - proj.

an injective resolution of  $X$  - resolution of the form

$\dots \rightarrow 0 \rightarrow I^0 \rightarrow I^1 \rightarrow \dots$   $I^i$  - injective.

Recall  $\mathcal{A}$  has enough projectives / injectives if  
 $\forall A \in \mathcal{A} \exists P \twoheadrightarrow A / A \twoheadrightarrow I$  with  
 $P$ -projective /  $I$ -injective.

If  $\mathcal{A}$  has enough projectives / injectives  $\Rightarrow$   
 $\Rightarrow$  every object has a projective / injective  
resolution.

Prop  $\text{Mod-}A$  &  $A\text{-Mod}$  have enough projectives.

Pf Free modules are projective, every module  
is a quotient of a free module.

Thm  $A\text{-Mod}$  &  $\text{Mod-}A$  have enough injectives.

Side remark: projectives are easier to deal with  
(usually), but in many important situations  
sheaves of  $\mathcal{A}$ s e.g., there is not enough of  
them.



Prop (Bayer's criterion)

$I \in A\text{-Mod}$  is injective  $\Leftrightarrow$  satisfies the UP  
for all  $I \hookrightarrow A$ ,  $I$ -ideal.

Pf  $\Rightarrow$  immediate

$$\begin{array}{ccc} \leftarrow & & \\ & 0 \rightarrow X & \xrightarrow{\iota} Y \\ & \uparrow f & \\ & I & \end{array}$$

consider the maximal submodule in  $Y$  containing  $X$   
for which the ext exists (Zorn).

WLOG we may assume that  $f$  does not extend  
to any larger submodule.

If  $Y \neq X \Rightarrow \exists y \in Y \setminus X$ .

Consider  $0 \rightarrow I \rightarrow X \oplus A \rightarrow X' \rightarrow 0$

$$\begin{array}{ccccc}
 I & \rightarrow & X & \xrightarrow{f} & \\
 \downarrow & \cong & \downarrow & \dashrightarrow & I \\
 A & \rightarrow & X' & \dashrightarrow & 
 \end{array}$$

$\downarrow$   
 $X \oplus (y)$

This square is both cartesian and cocartesian.  $\square$

Remark A very similar argument (though on ordinals) shows that in any Grothendieck category injective envelopes exist!

Application to A-Mod:

- 1) Check that  $\mathbb{Q}/\mathbb{Z}$  is injective in  $\mathcal{A}b$ !
- 2)  $F: \mathcal{A} \rightarrow \mathcal{B}$ ,  $G: \mathcal{B} \rightarrow \mathcal{A}$  are adjoint  $\neq$   
 $F$  is exact, then  $G$  sends injectives to injectives.

3) Consider functors

$$F: A\text{-Mod} \rightarrow (\text{Mod-}A)^{\circ} \quad G: (\text{Mod-}A)^{\circ} \rightarrow A\text{-Mod}$$

$$M \mapsto \text{Hom}_{\text{Ab}}(M, \mathbb{Q}/\mathbb{Z})$$

Check that they are adjoint & exact.  
Use that  $(\text{Mod-}A)^{\circ}$  has enough injectives:  
same as projectives in  $\text{Mod-}A$ .