

On channels with positive quantum zero-error capacity having vanishing n -shot capacity

M. E. Shirokov¹

Received: 22 January 2015 / Accepted: 28 April 2015 / Published online: 21 May 2015
© Springer Science+Business Media New York 2015

Abstract We show that unbounded number of channel uses may be necessary for perfect transmission of quantum information. For any n , we explicitly construct low-dimensional quantum channels (input dimension 4, Choi rank 2 or 4) whose quantum zero-error capacity is positive, but the corresponding n -shot capacity is zero. We give estimates for quantum zero-error capacity of such channels as a function of n and show that these channels can be chosen in any small vicinity (in the cb -norm) of a classical–quantum channel. Mathematically, this property means appearance of an ideal (noiseless) subchannel only in sufficiently large tensor power of a channel. Our approach (using special continuous deformation of a maximal commutative $*$ -subalgebra of M_4) also gives low-dimensional examples of the superactivation of 1-shot quantum zero-error capacity. Finally, we consider multi-dimensional construction which increases the estimate for quantum zero-error capacity of channels having vanishing n -shot capacity.

Keywords Pseudo-diagonal quantum channel · Error-correcting code · Noncommutative graph · Maximal commutative $*$ -algebra

1 Introduction

It is well known that the rate of information transmission over classical and quantum communication channels can be increased by simultaneous use of many copies of a

The research is funded by the Grant of Russian Science Foundation (Project No. 14-21-00162).

✉ M. E. Shirokov
msh@mi.ras.ru

¹ Steklov Mathematical Institute, Moscow, Russia

channel. It is this fact that implies necessity of regularization in definitions of different capacities of a channel [1, 2].

In this paper, we show that zero-error transmission of quantum information over a quantum channel may require unbounded number of channel uses. We prove by explicit construction that for any given n , there is a channel Φ_n such that

$$\bar{Q}_0(\Phi_n^{\otimes n}) = 0, \quad \text{but} \quad Q_0(\Phi_n) > 0, \quad (1)$$

where \bar{Q}_0 and Q_0 are, respectively, the 1-shot and the asymptotic quantum zero-error capacities defined in Sect. 2.

This effect is closely related to the recently discovered phenomenon of superactivation of quantum channel capacities [3–6]. Indeed, (1) is equivalent to existence of $m > n$ such that

$$\bar{Q}_0(\Phi_n) = \bar{Q}_0(\Phi_n^{\otimes 2}) = \dots = \bar{Q}_0(\Phi_n^{\otimes(m-1)}) = 0, \quad \text{but} \quad \bar{Q}_0(\Phi_n^{\otimes m}) > 0. \quad (2)$$

Mathematically, (2) means that all the channels $\Phi_n, \Phi_n^{\otimes 2}, \dots, \Phi_n^{\otimes(m-1)}$ have no ideal (noiseless) subchannels but the channel $\Phi_n^{\otimes m}$ has.

We show how for any given n to explicitly construct a pseudo-diagonal quantum channel Φ_n with the input dimension $d_A = 4$ and the Choi rank $d_E \geq 2$ satisfying (2) by determining its noncommutative graph. We also obtain the estimate for m as a function of n , which gives the lower bound for $Q_0(\Phi_n)$ in (1). This shows that

$$\sup_{\Phi} \{Q_0(\Phi) \mid \bar{Q}_0(\Phi^{\otimes n}) = 0\} \geq \frac{2 \ln(3/2)}{\pi n} \quad \forall n. \quad (3)$$

It is also observed that a channel Φ_n satisfying (1) and (2) can be obtained by arbitrarily small deformation (in the cb -norm) of a classical–quantum channel with $d_A = d_E = 4$.

The main problem in finding the channel Φ_n is to show nonexistence of error-correcting codes for the channel $\Phi_n^{\otimes n}$ (provided the existence of such codes is proved for $\Phi_n^{\otimes m}$). We solve this problem by using special continuous deformation of a maximal commutative $*$ -subalgebra of 4×4 matrices as the noncommutative graph of Φ_n and by noting that the Knill–Laflamme error-correcting conditions are violated for any maximal commutative $*$ -subalgebra with the positive dimension-independent gap (Lemma 3).

Our construction also gives low-dimensional examples of the superactivation of 1-shot quantum zero-error capacity. In particular, it gives an example of symmetric superactivation with $d_A = 4, d_E = 2$ (simplifying the example in [7]) and shows that such superactivation is possible for two channels with $d_A = d_E = 4$ if one of them is arbitrarily close (in the cb -norm) to a classical–quantum channel.

In the last section, we consider multi-dimensional generalization of our basic construction. It gives examples of channels which amplify the lower bound in (3) by the factor $\frac{\pi}{2 \ln 2} \approx 2.26$. Unfortunately, we did not managed to show that the value in the left side of (3) is $+\infty$ (as it is reasonable to conjecture). Estimation of this value remains an open question.

It should mention that the necessity of unbounded number of channel uses to see quantum ε -error capacity has been recently shown in [8].¹

2 Preliminaries

Let \mathcal{H} be a finite-dimensional Hilbert space, $\mathfrak{B}(\mathcal{H})$ —the algebra of all linear operators in \mathcal{H} , $\mathfrak{S}(\mathcal{H})$ —the closed convex subset of $\mathfrak{B}(\mathcal{H})$ consisting of positive operators with unit trace called *quantum states* [1, 2]. We can identify \mathcal{H} and $\mathfrak{B}(\mathcal{H})$, respectively, with the unitary space \mathbb{C}^n and the algebra \mathfrak{M}_n of all $n \times n$ matrices, where $n = \dim \mathcal{H}$.

Denote by I_n and Id_n the unit operator in the space \mathbb{C}^n and the identity transformation of the algebra \mathfrak{M}_n correspondingly.

For any matrix $A \in \mathfrak{M}_n$ denote by Υ_A the operator of Schur multiplication by A in \mathfrak{M}_n (also called the Hadamard multiplication). Its *cb*-norm will be denoted $\|\Upsilon_A\|_{\text{cb}}$. It coincides with the operator norm of Υ_A and is also called the Schur (or Hadamard) multiplier norm of A (see [9, 10] and the references therein).

Let $\Phi : \mathfrak{S}(\mathcal{H}_A) \rightarrow \mathfrak{S}(\mathcal{H}_B)$ be a quantum channel, i.e., a completely positive trace-preserving linear map [1, 2]. Stinespring's theorem implies the existence of a Hilbert space \mathcal{H}_E and of an isometry $V : \mathcal{H}_A \rightarrow \mathcal{H}_B \otimes \mathcal{H}_E$ such that

$$\Phi(\rho) = \text{Tr}_{\mathcal{H}_E} V \rho V^*, \quad \rho \in \mathfrak{S}(\mathcal{H}_A).$$

The minimal dimension of \mathcal{H}_E is called Choi rank of Φ and denoted d_E .

The quantum channel

$$\mathfrak{S}(\mathcal{H}_A) \ni \rho \mapsto \widehat{\Phi}(\rho) = \text{Tr}_{\mathcal{H}_B} V \rho V^* \in \mathfrak{S}(\mathcal{H}_E)$$

is called *complementary* to the channel Φ [1, 11]. The complementary channel is defined uniquely up to isometrical equivalence [11, the Appendix].

The 1-shot quantum zero-error capacity $\bar{Q}_0(\Phi)$ of a channel Φ is defined as $\sup_{\mathcal{H} \in q_0(\Phi)} \log_2 \dim \mathcal{H}$, where $q_0(\Phi)$ is the set of all subspaces \mathcal{H}_0 of \mathcal{H}_A on which the channel Φ is perfectly reversible (in the sense that there is a channel Θ such that $\Theta(\Phi(\rho)) = \rho$ for all states ρ supported by \mathcal{H}_0). Any subspace $\mathcal{H}_0 \in q_0(\Phi)$ is called *error-correcting code* for the channel Φ [1, 13].

The (asymptotic) quantum zero-error capacity is defined by regularization: $Q_0(\Phi) = \sup_n n^{-1} \bar{Q}_0(\Phi^{\otimes n})$ [6, 12, 13].

It is well known that a channel Φ is perfectly reversible on a subspace \mathcal{H}_0 if and only if the restriction of the complementary channel $\widehat{\Phi}$ to the subset $\mathfrak{S}(\mathcal{H}_0)$ is completely depolarizing, i.e., $\widehat{\Phi}(\rho_1) = \widehat{\Phi}(\rho_2)$ for all states ρ_1 and ρ_2 supported by \mathcal{H}_0 [1, Ch.10]. It follows that the 1-shot quantum zero-error capacity of a channel Φ is completely determined by the set $\mathcal{G}(\Phi) \doteq \widehat{\Phi}^*(\mathfrak{B}(\mathcal{H}_E))$ called *noncommutative graph* of Φ [13]. In particular, the Knill–Laflamme error-correcting condition (cf. [14]) implies the following lemma.

¹ It is surprising that this result and the preliminary arXiv version of the present paper appeared simultaneously.

Lemma 1 A set $\{\varphi_k\}_{k=1}^d$ of unit orthogonal vectors in \mathcal{H}_A is a basis of error-correcting code for a channel $\Phi : \mathfrak{S}(\mathcal{H}_A) \rightarrow \mathfrak{S}(\mathcal{H}_B)$ if and only if

$$\langle \varphi_l | A | \varphi_k \rangle = 0 \quad \text{and} \quad \langle \varphi_l | A | \varphi_l \rangle = \langle \varphi_k | A | \varphi_k \rangle \quad \forall A \in \mathfrak{L}, \quad \forall k \neq l, \quad (4)$$

where \mathfrak{L} is any subset of $\mathfrak{B}(\mathcal{H}_A)$ such that $\text{lin} \mathfrak{L} = \mathcal{G}(\Phi)$.

This lemma shows that $\bar{Q}_0(\Phi) \geq \log_2 d$ if and only if there exists a set $\{\varphi_k\}_{k=1}^d$ of unit vectors in \mathcal{H}_A satisfying condition (4).

Remark 1 Since a subspace \mathfrak{L} of the algebra \mathfrak{M}_n of $n \times n$ matrices is a noncommutative graph of a particular channel if and only if

$$\mathfrak{L} \text{ is symmetric } (\mathfrak{L} = \mathfrak{L}^*) \text{ and contains the unit matrix} \quad (5)$$

(see Lemma 2 in [6] or Proposition 2 in [15]), Lemma 1 shows that one can “construct” a channel Φ with $\dim \mathcal{H}_A = n$ having positive (correspondingly, zero) 1-shot quantum zero-error capacity by taking a subspace $\mathfrak{L} \subset \mathfrak{M}_n$ satisfying (5) for which the following condition is valid (correspondingly, not valid)

$$\exists \varphi, \psi \in [\mathbb{C}^n]_1 \text{ such that } \langle \psi | A | \varphi \rangle = 0 \text{ and } \langle \varphi | A | \varphi \rangle = \langle \psi | A | \psi \rangle \quad \forall A \in \mathfrak{L}, \quad (6)$$

where $[\mathbb{C}^n]_1$ is the unit sphere of \mathbb{C}^n . \square

We will use the following two notions.

Definition 1 [1] A channel $\Phi : \mathfrak{S}(\mathcal{H}_A) \rightarrow \mathfrak{S}(\mathcal{H}_B)$ is called *classical–quantum* if it has the representation

$$\Phi(\rho) = \sum_k \langle k | \rho | k \rangle \sigma_k,$$

where $\{|k\rangle\}$ is an orthonormal basis in \mathcal{H}_A and $\{\sigma_k\}$ is a collection of states in $\mathfrak{S}(\mathcal{H}_B)$.

Definition 2 [16] A channel $\Phi : \mathfrak{S}(\mathcal{H}_A) \rightarrow \mathfrak{S}(\mathcal{H}_B)$ is called *pseudo-diagonal* if it has the representation

$$\Phi(\rho) = \sum_{i,j} c_{ij} \langle \psi_i | \rho | \psi_j \rangle |i\rangle \langle j|,$$

where $\{c_{ij}\}$ is a Gram matrix of a collection of unit vectors, $\{|\psi_i\rangle\}$ is a collection of vectors in \mathcal{H}_A such that $\sum_i |\psi_i\rangle \langle \psi_i| = I_{\mathcal{H}_A}$ – the unit operator in \mathcal{H}_A and $\{|i\rangle\}$ is an orthonormal basis in \mathcal{H}_B .

Pseudo-diagonal channels are complementary to entanglement-breaking channels and vice versa [11, 16].

3 Basic example

For any given $\theta \in T \doteq (-\pi, \pi]$, consider the subspace

$$\mathfrak{L}_\theta = \left\{ M = \begin{bmatrix} a & b & \gamma c & d \\ b & a & d & \bar{\gamma} c \\ \bar{\gamma} c & d & a & b \\ d & \gamma c & b & a \end{bmatrix}, \quad a, b, c, d \in \mathbb{C}, \quad \gamma = \exp\left(\frac{i}{2}\theta\right) \right\} \quad (7)$$

of \mathfrak{M}_4 . This subspace satisfies condition (5) and has the following property

$$A = W_4^* A W_4 \quad \forall A \in \mathfrak{L}_\theta, \quad \text{where} \quad W_4 = \begin{bmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}. \quad (8)$$

Remark 2 The subspace \mathfrak{L}_θ can be considered as a deformation of the maximal commutative $*$ -algebra \mathfrak{L}_0 . To clarify the form of this deformation, note that the family of subspaces \mathfrak{L}_θ is unitary equivalent to the family of the subspaces

$$\mathfrak{L}_\theta^s = \left\{ M = \begin{bmatrix} a & 0 & 0 & 0 \\ 0 & b & 0 & 0 \\ 0 & 0 & c & 0 \\ 0 & 0 & 0 & d \end{bmatrix} + \frac{1}{4}(d + c - b - a) T_\theta, \quad a, b, c, d \in \mathbb{C} \right\} \quad (9)$$

where

$$T_\theta = \begin{bmatrix} u & 0 & 0 & v \\ 0 & u & v & 0 \\ 0 & -v & -u & 0 \\ -v & 0 & 0 & -u \end{bmatrix}, \quad u = 1 - \Re \gamma, \quad v = i \Im \gamma.$$

Indeed, by representing the matrix M in (7) as $M = A + cB$, where

$$A = \begin{bmatrix} a & b & c & d \\ b & a & d & c \\ c & d & a & b \\ d & c & b & a \end{bmatrix}, \quad B = \begin{bmatrix} 0 & 0 & \tau & 0 \\ 0 & 0 & 0 & \bar{\tau} \\ \bar{\tau} & 0 & 0 & 0 \\ 0 & \tau & 0 & 0 \end{bmatrix}, \quad \tau = \gamma - 1,$$

is easy to see that $S^{-1}BS = T_\theta$ and

$$S^{-1}AS = \begin{bmatrix} \tilde{a} & 0 & 0 & 0 \\ 0 & \tilde{b} & 0 & 0 \\ 0 & 0 & \tilde{c} & 0 \\ 0 & 0 & 0 & \tilde{d} \end{bmatrix}, \quad \text{where} \quad S = \frac{1}{2} \begin{bmatrix} 1 & 1 & 1 & 1 \\ -1 & 1 & -1 & 1 \\ -1 & -1 & 1 & 1 \\ 1 & -1 & -1 & 1 \end{bmatrix},$$

$$\tilde{a} = a - b - c + d, \quad \tilde{b} = a + b - c - d, \quad \tilde{c} = a - b + c - d, \quad \tilde{d} = a + b + c + d.$$

Denote by $\widehat{\mathfrak{L}}_\theta$ the set of all channels whose noncommutative graph coincides with \mathfrak{L}_θ . Since $\dim \mathfrak{L}_\theta = 4$, Proposition 2 in [15] shows the set $\widehat{\mathfrak{L}}_\theta$ contains infinitely many different channels with $d_A \doteq \dim \mathcal{H}_A = 4$ and $d_E \geq 2$. It is essential that one can choose families $\{\Phi_\theta\} \subset \{\widehat{\mathfrak{L}}_\theta\}$ continuous with respect to θ .

Lemma 2 1) *There is a family $\{\Phi_\theta^1\}$ of pseudo-diagonal channels (see Def. 2) with $d_E = 2$ such that $\Phi_\theta^1 \in \widehat{\mathfrak{L}}_\theta$ for each θ .*
 2) *There is a family $\{\Phi_\theta^2\}$ of pseudo-diagonal channels with $d_E = 4$ such that $\Phi_\theta^2 \in \widehat{\mathfrak{L}}_\theta$ for each θ and Φ_0^2 is a classical–quantum channel (see Def. 1).
 The families $\{\Phi_\theta^1\}$ and $\{\Phi_\theta^2\}$ can be chosen continuous in the following sense:*

$$\Phi_\theta^k(\rho) = \text{Tr}_{\mathcal{H}_E^k} V_\theta^k \rho \left[V_\theta^k \right]^*, \quad \rho \in \mathfrak{S}(\mathcal{H}_A), \quad k = 1, 2, \quad (10)$$

where V_θ^1, V_θ^2 are continuous families of isometries, $\mathcal{H}_E^1 = \mathbb{C}^2$, $\mathcal{H}_E^2 = \mathbb{C}^4$.²

Lemma 2 is proved in the Appendix, where representations (10) are constructed explicitly by using the unitary equivalence of \mathfrak{L}_θ and \mathfrak{L}_θ^s .

Theorem 1 *Let Φ_θ be a channel in $\widehat{\mathfrak{L}}_\theta$ and $n \in \mathbb{N}$ be arbitrary.*

- A) $\bar{Q}_0(\Phi_\theta) > 0$ if and only if $\theta = \pi$ and $\bar{Q}_0(\Phi_\pi) = 1$.
 B) *If $\theta_1 + \dots + \theta_n = \pi \pmod{2\pi}$, then $\bar{Q}_0(\Phi_{\theta_1} \otimes \dots \otimes \Phi_{\theta_n}) > 0$ and there exist 2^n mutually orthogonal 2-D error-correcting codes for the channel $\Phi_{\theta_1} \otimes \dots \otimes \Phi_{\theta_n}$. For each binary n -tuple (x_1, \dots, x_n) , the corresponding error-correcting code is spanned by the images of the vectors*

$$|\varphi\rangle = \frac{1}{\sqrt{2}} [|1\dots 1\rangle + i |2\dots 2\rangle], \quad |\psi\rangle = \frac{1}{\sqrt{2}} [|3\dots 3\rangle + i |4\dots 4\rangle], \quad (11)$$

under the unitary transformation $U_{x_1} \otimes \dots \otimes U_{x_n}$, where $\{|1\rangle, \dots, |4\rangle\}$ is the canonical basis in \mathbb{C}^4 , $U_0 = I_4$ and $U_1 = W_4$ (defined in (8)).

- C) *If $|\theta_1| + \dots + |\theta_n| \leq 2 \ln(3/2)$, then $\bar{Q}_0(\Phi_{\theta_1} \otimes \dots \otimes \Phi_{\theta_n}) = 0$.*

Remark 3 It is easy to show that $\bar{Q}_0(\Phi_\theta^{\otimes n}) = \bar{Q}_0(\Phi_{-\theta}^{\otimes n})$ and that the set of all θ such that $\bar{Q}_0(\Phi_\theta^{\otimes n}) = 0$ is open. Hence, for each n , there is $\varepsilon_n > 0$ such that $\bar{Q}_0(\Phi_\theta^{\otimes n}) = 0$ if $|\theta| < \varepsilon_n$ and $\bar{Q}_0(\Phi_{\pm\varepsilon_n}^{\otimes n}) > 0$. Theorem 1 shows that $\varepsilon_1 = \pi$ and $2 \ln(3/2)/n < \varepsilon_n \leq \pi/n$ for $n > 1$. Since assertion C is proved by using quite coarse estimates, one can conjecture that $\varepsilon_n = \pi/n$ for $n > 1$. There exist some arguments confirming validity of this conjecture for $n = 2$.

Remark 4 Assertion B of Theorem 1 can be strengthened as follows:

B') *If $\theta_1 + \dots + \theta_n = \pi \pmod{2\pi}$, then there exist 2^n mutually orthogonal projectors $P_{\bar{x}}$ of rank 2 indexed by a binary n -tuple $\bar{x} = (x_1, \dots, x_n)$ such that*

$$P_{\bar{x}} A P_{\bar{x}} = \lambda(A) P_{\bar{x}} \quad \forall A \in \mathfrak{L}_{\theta_1} \otimes \dots \otimes \mathfrak{L}_{\theta_n},$$

² This implies continuity of these families in the cb -norm [17].

where $\lambda(A) \in \mathbb{C}$ does not depend on \bar{x} . $P_{\bar{x}}$ is the projector on the subspace $U_{x_1} \otimes \dots \otimes U_{x_n}(\mathcal{H}_0)$, where \mathcal{H}_0 is the linear hull of vectors (11).

This follows from the proof of Theorem 1 presented below.

Theorem 1 implies the main result of this paper.

Corollary 1 *Let n be arbitrary and m be a natural number such that $\theta_* = \pi/m \leq 2 \ln(3/2)/n$. Then*

$$\bar{Q}_0(\Phi_{\theta_*}^{\otimes n}) = 0 \quad \text{but} \quad \bar{Q}_0(\Phi_{\theta_*}^{\otimes m}) \geq 1 \quad \text{and hence} \quad Q_0(\Phi_{\theta_*}) \geq 1/m. \quad (12)$$

There exist 2^m mutually orthogonal 2-D error-correcting codes for the channel $\Phi_{\theta_*}^{\otimes m}$.

Relation (12) means that it is not possible to transmit any quantum information with no errors by using $\leq n$ copies of the channel Φ_{θ_*} , but such transmission is possible if the number of copies is $\geq m$.

Remark 5 In (12), one can take $\Phi_{\theta_*} = \Phi_{\theta_*}^1$ —a channel from the family described in the first part of Lemma 2. So, Corollary 1 shows that for any n , there exists a channel Φ_n with $d_A = 4$ and $d_E = 2$ such that $\bar{Q}_0(\Phi_n^{\otimes n}) = 0$ and

$$Q_0(\Phi_n) \geq \left(\left\lfloor \frac{\pi n}{2 \ln(3/2)} \right\rfloor + 1 \right)^{-1} = \frac{2 \ln(3/2)}{\pi n} + o(1/n), \quad n \rightarrow +\infty,$$

where $[x]$ is the integer part of x .

It is natural to ask about the maximal value of quantum zero-error capacity of a channel with given input dimension having vanishing n -shot capacity, i.e., about the value

$$S_d(n) \doteq \sup_{\Phi: d_A=d} \{Q_0(\Phi) \mid \bar{Q}_0(\Phi^{\otimes n}) = 0\}, \quad (13)$$

where the supremum is over all quantum channels with $d_A \doteq \dim \mathcal{H}_A = d$. We may also consider the value

$$S_*(n) \doteq \sup_d S_d(n) = \lim_{d \rightarrow +\infty} S_d(n) \leq +\infty. \quad (14)$$

The sequences $\{S_d(n)\}_n$ and $\{S_*(n)\}_n$ are nonincreasing, and the first of them is bounded by $\log_2 d$. Theorem 2 in [7] shows that

$$S_{2d}(1) \geq \frac{\log_2 d}{2} \quad \text{and hence} \quad S_*(1) = +\infty.$$

It seems reasonable to conjecture that $S_*(n) = +\infty$ for all n . A possible way to prove this conjecture is discussed at the end of Sect. 4.

It follows from the superadditivity of quantum zero-error capacity that

$$S_{d^k}(n) \geq kS_d(nk) \quad \text{and hence} \quad S_*(n) \geq kS_*(nk) \quad \text{for any } k, n. \quad (15)$$

These relations show that the assumption $S_*(n_0) < +\infty$ for some n_0 implies

$$S_d(n) = O(1/n) \quad \text{for each } d \quad \text{and} \quad S_*(n) = O(1/n) \quad \text{if } n \geq n_0.$$

By Corollary 1, we have

$$S_4(n) \geq \left(\left\lfloor \frac{\pi n}{2 \ln(3/2)} \right\rfloor + 1 \right)^{-1} = \frac{2 \ln(3/2)}{\pi n} + o(1/n), \quad \forall n. \quad (16)$$

This and (15) imply the estimation

$$S_{4^k}(n) \geq k \frac{2 \ln(3/2)}{\pi kn} + o(1/(kn)) = \frac{2 \ln(3/2)}{\pi n} + o(1/(kn)),$$

which shows that

$$S_*(n) \geq \frac{2 \ln(3/2)}{\pi n} \quad \forall n. \quad (17)$$

In Sect. 4, we will improve these lower bounds by considering the multi-dimensional generalization of the above construction.

Remark 6 Since the parameter θ_* in Corollary 1 can be taken arbitrarily close to zero, the second part of Lemma 2 shows that the channel Φ_{θ_*} , for which $\bar{Q}_0(\Phi_{\theta_*}^{\otimes n}) = 0$ and $Q_0(\Phi_{\theta_*}) > 0$, can be chosen in any small vicinity (in the cb -norm) of the classical-quantum channel Φ_0^2 .

Theorem 1 also gives examples of the superactivation of 1-shot quantum zero-error capacity.

Corollary 2 *If $\theta \neq 0, \pi$, then the following superactivation property*

$$\bar{Q}_0(\Phi_\theta) = \bar{Q}_0(\Phi_{\pi-\theta}) = 0 \quad \text{and} \quad \bar{Q}_0(\Phi_\theta \otimes \Phi_{\pi-\theta}) > 0$$

holds for any channels $\Phi_\theta \in \widehat{\mathfrak{L}}_\theta$ and $\Phi_{\pi-\theta} \in \widehat{\mathfrak{L}}_{\pi-\theta}$. For any $\theta \in \mathbb{T}$, there exist 4 mutually orthogonal 2-D error-correcting codes for the channel $\Phi_\theta \otimes \Phi_{\pi-\theta}$, one of them is spanned by the vectors

$$|\varphi\rangle = \frac{1}{\sqrt{2}} [|11\rangle + i |22\rangle], \quad |\psi\rangle = \frac{1}{\sqrt{2}} [|33\rangle + i |44\rangle], \quad (18)$$

others are the images of this subspace under the unitary transformations $I_4 \otimes W_4$, $W_4 \otimes I_4$ and $W_4 \otimes W_4$ (the operator W_4 is defined in (8)).

Remark 7 Corollary 2 shows that the channel $\Phi_{\pi/2}^1$ (taken from the first part of Lemma 2) is an example of the symmetric superactivation of 1-shot quantum zero-error capacity with Choi rank 2.³

By taking the family $\{\Phi_\theta^2\}$ from the second part of Lemma 2 and tending θ to zero, we see from Corollary 2 that the superactivation of 1-shot quantum zero-error capacity may hold for two channels with $d_A = d_E = 4$ if one of them is arbitrarily close (in the cb-norm) to a classical–quantum channel.

Note that the entangled subspace spanned by the vectors (18) is an error-correcting code for the channel $\Phi_0^2 \otimes \Phi_\pi^2$ (and hence for the channel $\Phi_0^2 \otimes \text{Id}_4$) despite the fact that Φ_0^2 is a classical–quantum channel.

Proof of Theorem 1 A) It is easy to verify that the subspace \mathfrak{L}_π satisfies condition (6) with the vectors $|\varphi\rangle = [1, i, 0, 0]^\top$, $|\psi\rangle = [0, 0, 1, i]^\top$.

To prove that $\bar{Q}_0(\Phi_\theta) = 0$ for all $\theta \neq \pi$, it suffices to show that condition (6) is not valid for the subspace \mathfrak{L}_θ^s defined in (9) if $\theta \neq \pi$ (i.e., $\gamma \neq i$).

Assume the existence of unit vectors $|\varphi\rangle = [x_1, x_2, x_3, x_4]^\top$ and $|\psi\rangle = [y_1, y_2, y_3, y_4]^\top$ in \mathbb{C}^4 such that

$$\langle\psi|M|\varphi\rangle = 0 \quad \text{and} \quad \langle\psi|M|\psi\rangle = \langle\varphi|M|\varphi\rangle \quad \text{for all } M \in \mathfrak{L}_\theta^s. \quad (19)$$

Since condition (19) is invariant under the rotation

$$|\varphi\rangle \mapsto p|\varphi\rangle - q|\psi\rangle, \quad |\psi\rangle \mapsto \bar{q}|\varphi\rangle + \bar{p}|\psi\rangle, \quad |p|^2 + |q|^2 = 1,$$

we may consider that $y_1 = 0$.

By taking successively $(a = -1, b = c = d = 0)$, $(b = -1, a = c = d = 0)$, $(c = 1, a = b = d = 0)$ and $(d = 1, a = b = c = 0)$, we obtain from (19) the following equations

$$\begin{aligned} \bar{y}_1 x_1 &= \bar{y}_2 x_2 = -\bar{y}_3 x_3 = -\bar{y}_4 x_4 = \frac{1}{4} \langle\psi|T_\theta|\varphi\rangle, \\ |x_1|^2 - |y_1|^2 &= |x_2|^2 - |y_2|^2 = |y_3|^2 - |x_3|^2 = |y_4|^2 - |x_4|^2 \\ &= \frac{1}{4} (\langle\varphi|T_\theta|\varphi\rangle - \langle\psi|T_\theta|\psi\rangle). \end{aligned}$$

Since $y_1 = 0$ and $\|\varphi\| = \|\psi\| = 1$, the above equations imply

$$y_1 = y_2 = x_3 = x_4 = 0$$

and

$$|x_1|^2 = |x_2|^2 = |y_3|^2 = |y_4|^2 = \frac{1}{4} (\langle\varphi|T_\theta|\varphi\rangle - \langle\psi|T_\theta|\psi\rangle) = 1/2. \quad (20)$$

So, $|\varphi\rangle = [x_1, x_2, 0, 0]^\top$ and $|\psi\rangle = [0, 0, y_3, y_4]^\top$, where $[x_1, x_2]^\top$ and $[y_3, y_4]^\top$ are unit vectors in \mathbb{C}^2 . It follows from (20) that

³ This strengthens the result in [7], where a similar example with Choi rank 3 and the same input dimension was constructed.

$$2 = \left\langle x_1 \middle| \begin{array}{cc} u & 0 \\ 0 & u \end{array} \middle| x_1 \right\rangle - \left\langle y_3 \middle| \begin{array}{cc} -u & 0 \\ 0 & -u \end{array} \middle| y_3 \right\rangle = 2u,$$

which can be valid only if $\gamma = i$, i.e., $\theta = \pi$.

The above arguments also show that $\bar{Q}_0(\Phi_\pi) = 1$, since the assumption $\bar{Q}_0(\Phi_\pi) > 1$ implies, by Lemma 1, the existence of orthogonal unit vectors ϕ_1, ϕ_2, ϕ_3 such that condition (19) with $\varphi = \phi_i, \psi = \phi_j$ is valid for all $i \neq j$.

B) Let $M_1 \in \mathfrak{L}_{\theta_1}, \dots, M_n \in \mathfrak{L}_{\theta_n}$ be arbitrary and $X = M_1 \otimes \dots \otimes M_n$. To prove that the linear hull \mathcal{H}_0 of vectors (11) is an error-correcting code for the channel $\Phi_{\theta_1} \otimes \dots \otimes \Phi_{\theta_n}$, it suffices, by Lemma 1, to show that

$$\langle \psi | X | \varphi \rangle = 0 \quad \text{and} \quad \langle \psi | X | \psi \rangle = \langle \varphi | X | \varphi \rangle. \quad (21)$$

We have

$$\begin{aligned} 2\langle \psi | X | \varphi \rangle &= \langle 3 \dots 3 | X | 1 \dots 1 \rangle + i\langle 3 \dots 3 | X | 2 \dots 2 \rangle - i\langle 4 \dots 4 | X | 1 \dots 1 \rangle \\ &\quad + \langle 4 \dots 4 | X | 2 \dots 2 \rangle = c_1 \dots c_n (\bar{\gamma}_1 \dots \bar{\gamma}_n + \gamma_1 \dots \gamma_n) \\ &\quad + d_1 \dots d_n (i - i) = 0, \end{aligned}$$

since $\gamma_1 \dots \gamma_n = \pm i$,

$$\begin{aligned} 2\langle \varphi | X | \varphi \rangle &= \langle 1 \dots 1 | X | 1 \dots 1 \rangle + i\langle 1 \dots 1 | X | 2 \dots 2 \rangle - i\langle 2 \dots 2 | X | 1 \dots 1 \rangle \\ &\quad + \langle 2 \dots 2 | X | 2 \dots 2 \rangle = a_1 \dots a_n (1 + 1) + b_1 \dots b_n (i - i) = 2a_1 \dots a_n \end{aligned}$$

and

$$\begin{aligned} 2\langle \psi | X | \psi \rangle &= \langle 3 \dots 3 | X | 3 \dots 3 \rangle + i\langle 3 \dots 3 | X | 4 \dots 4 \rangle - i\langle 4 \dots 4 | X | 3 \dots 3 \rangle \\ &\quad + \langle 4 \dots 4 | X | 4 \dots 4 \rangle = a_1 \dots a_n (1 + 1) + b_1 \dots b_n (i - i) = 2a_1 \dots a_n. \end{aligned}$$

So, the both equalities in (21) are valid.

To prove that the subspace $U_{\bar{x}}(\mathcal{H}_0)$, where $U_{\bar{x}} = U_{x_1} \otimes \dots \otimes U_{x_n}$, is an error-correcting code for the channel $\Phi_{\theta_1} \otimes \dots \otimes \Phi_{\theta_n}$, it suffices to note that (8) implies that $U_{\bar{x}}^* A U_{\bar{x}} = A$ for all $A \in \mathfrak{L}_{\theta_1} \otimes \dots \otimes \mathfrak{L}_{\theta_n}$.

C) To show that $\bar{Q}_0(\bigotimes_{k=1}^n \Phi_{\theta_k}) = 0$ if $\sum_{k=1}^n |\theta_k| \leq 2 \ln(3/2)$ note that $\mathfrak{L}_\theta = \Upsilon_{D(\theta)}(\mathfrak{L}_0)$ and $\bigotimes_{k=1}^n \mathfrak{L}_{\theta_k} = \bigotimes_{k=1}^n \Upsilon_{D(\theta_k)}(\mathfrak{L}_0^{\otimes n})$, where $\Upsilon_{D(\theta)}$ is the Schur multiplication by the matrix

$$D(\theta) = \begin{bmatrix} 1 & 1 & \gamma & 1 \\ 1 & 1 & 1 & \bar{\gamma} \\ \bar{\gamma} & 1 & 1 & 1 \\ 1 & \gamma & 1 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{bmatrix} + \begin{bmatrix} 0 & 0 & \tau & 0 \\ 0 & 0 & 0 & \bar{\tau} \\ \bar{\tau} & 0 & 0 & 0 \\ 0 & \tau & 0 & 0 \end{bmatrix}, \quad \tau = \gamma - 1. \quad (22)$$

By using (22) and Theorem 8.7 in [9]⁴, it is easy to show that

$$x_k \doteq \|\Upsilon_{D(\theta_k)} - \text{Id}_4\|_{\text{cb}} \leq |\tau_k| = |1 - \gamma_k| = \left|1 - \exp\left(\frac{i}{2}\theta_k\right)\right| \leq \frac{1}{2}|\theta_k|. \quad (23)$$

Let $\Delta_n \doteq \|\bigotimes_{k=1}^n \Upsilon_{D(\theta_k)} - \text{Id}_{4^n}\|_{\text{cb}}$. By using multiplicativity of the cb -norm and (23), we obtain

$$\Delta_n \leq x_n \prod_{k=1}^{n-1} (1 + x_k) + \Delta_{n-1} \leq \prod_{k=1}^n (1 + x_k) - 1 \leq \prod_{k=1}^n \left(1 + \frac{1}{2}|\theta_k|\right) - 1. \quad (24)$$

Assume that $\bar{Q}_0(\bigotimes_{k=1}^n \Phi_{\theta_k}) > 0$. Then, Lemma 1 implies existence of unit vectors φ and ψ in $\mathcal{H}_A^{\otimes n} = \mathbb{C}^{4^n}$ such that

$$\langle \psi | \Psi(A) | \varphi \rangle = 0 \quad \text{and} \quad \langle \varphi | \Psi(A) | \varphi \rangle = \langle \psi | \Psi(A) | \psi \rangle \quad \forall A \in \mathfrak{L}_0^{\otimes n},$$

where $\Psi = \bigotimes_{k=1}^n \Upsilon_{D(\theta_k)}$. Hence, for any A in the unit ball of $\mathfrak{L}_0^{\otimes n}$, we have

$$|\langle \psi | A | \varphi \rangle| \leq \Delta_n \quad \text{and} \quad |\langle \varphi | A | \varphi \rangle - \langle \psi | A | \psi \rangle| \leq 2\Delta_n.$$

By using (24) and the inequality $x \geq \ln(1+x)$, it is easy to see that the assumption $\sum_{k=1}^n |\theta_k| \leq 2 \ln(3/2)$ implies $\Delta_n \leq 1/2$. So, the above relations can not be valid by the below Lemma 3, since $\mathfrak{L}_0^{\otimes n}$ is a maximal commutative $*$ -subalgebra of \mathfrak{M}_{4^n} . \square

Lemma 3 *Let \mathfrak{A} be a maximal commutative $*$ -subalgebra of \mathfrak{M}_n . Then,*

$$\text{either } 2 \sup_{A \in \mathfrak{A}_1} |\langle \psi | A | \varphi \rangle| > 1 \quad \text{or} \quad \sup_{A \in \mathfrak{A}_1} |\langle \varphi | A | \varphi \rangle - \langle \psi | A | \psi \rangle| > 1$$

for any two unit vectors φ and ψ in \mathbb{C}^n , where \mathfrak{A}_1 is the unit ball of \mathfrak{A} .

Proof Let $\{x_i\}_{i=1}^n$ and $\{y_i\}_{i=1}^n$ be the coordinates of φ and ψ in the basis in which the algebra \mathfrak{A} consists of diagonal matrices. Then,

$$\sup_{A \in \mathfrak{A}_1} |\langle \psi | A | \varphi \rangle| = \sum_{i=1}^n |x_i| |y_i|, \quad \sup_{A \in \mathfrak{A}_1} |\langle \varphi | A | \varphi \rangle - \langle \psi | A | \psi \rangle| = \sum_{i=1}^n \left| |x_i|^2 - |y_i|^2 \right|.$$

Let $d_i = |y_i| - |x_i|$. Assume that

$$2 \sum_{i=1}^n |x_i| |y_i| \leq 1 \quad \text{and} \quad \sum_{i=1}^n \left| |x_i|^2 - |y_i|^2 \right| \leq 1.$$

⁴ This theorem states that $\|\Upsilon_A\|_{\text{cb}} \leq 1$ if and only if $A = [\langle \varphi_i | \psi_j \rangle]$ for some collections $\{\varphi_i\}$ and $\{\psi_j\}$ of vectors from the unit ball of some Hilbert space.

Since $\sum_{i=1}^n |x_i|^2 = \sum_{i=1}^n |y_i|^2 = 1$, the first of these inequalities implies

$$\left| \sum_{i=1}^n d_i |x_i| \right| \geq 1/2 \quad \text{and} \quad \left| \sum_{i=1}^n d_i |y_i| \right| \geq 1/2.$$

Hence,

$$\sum_{i=1}^n \left| |x_i|^2 - |y_i|^2 \right| = \sum_{i=1}^n |d_i| [|x_i| + |y_i|] > \left| \sum_{i=1}^n d_i |x_i| \right| + \left| \sum_{i=1}^n d_i |y_i| \right| \geq 1,$$

where the strict inequality follows from the existence of negative and positive numbers in the set $\{d_i\}_{i=1}^n$. This contradicts to the above assumption. \square

4 Multi-dimensional generalization

Note that

$$\mathfrak{L}_0 = \mathfrak{A}_2^{\otimes 2}, \quad \text{where} \quad \mathfrak{A}_2 = \left\{ \begin{bmatrix} a & b \\ b & a \end{bmatrix}, \quad a, b \in \mathbb{C} \right\},$$

and that \mathfrak{L}_θ is the image of \mathfrak{L}_0 under the Schur multiplication by matrix (22). So, the above construction can be generalized by considering the corresponding deformation of the maximal commutative $*$ -subalgebra $\mathfrak{L}_0^p = \mathfrak{A}_2^{\otimes p}$ of \mathfrak{M}_{2^p} for $p > 2$. The algebra \mathfrak{L}_0^p can be described recursively as follows:

$$\mathfrak{L}_0^p = \left\{ \begin{bmatrix} A & B \\ B & A \end{bmatrix}, \quad A, B \in \mathfrak{L}_0^{p-1} \right\}, \quad \mathfrak{L}_0^1 = \mathfrak{A}_2.$$

Let $p > 2$ and $\theta \in \mathbb{T} \doteq (-\pi, \pi]$ be arbitrary, $\gamma = \exp(\frac{i}{2}\theta)$. Let $D(\theta)$ be the $2^p \times 2^p$ matrix described as $2^{p-1} \times 2^{p-1}$ matrix $[A_{ij}]$ consisting of the blocks

$$A_{ii} = \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix} \quad \forall i, \quad A_{ij} = \begin{bmatrix} \gamma & 1 \\ 1 & \bar{\gamma} \end{bmatrix} \quad \text{if } i < j \quad \text{and} \quad A_{ij} = \begin{bmatrix} \bar{\gamma} & 1 \\ 1 & \gamma \end{bmatrix} \quad \text{if } i > j.$$

Consider the 2^p -D subspace $\mathfrak{L}_\theta^p = \Upsilon_{D(\theta)}(\mathfrak{L}_0^p)$ of \mathfrak{M}_{2^p} (where $\Upsilon_{D(\theta)}$ is the Schur multiplication by the matrix $D(\theta)$), see Sect. 2). This subspace satisfies condition (5) and has the following property

$$A = W_{2^p}^* A W_{2^p} \quad \forall A \in \mathfrak{L}_\theta^p, \quad (25)$$

where W_{2^p} is the $2^p \times 2^p$ matrix having “1” on the main skew-diagonal and “0” on the other places. To prove (25), it suffices to show that it holds for the algebra $\mathfrak{L}_0^p = \mathfrak{A}_2^{\otimes p}$ (by using $W_{2^p} = W_2^{\otimes p}$) and to note that the map $\Upsilon_{D(\theta)}$ commutes with the transformation $A \mapsto W_{2^p}^* A W_{2^p}$.

Denote by $\widehat{\mathfrak{L}}_\theta^p$ the set of all quantum channels whose noncommutative graph coincides with \mathfrak{L}_θ^p . By Proposition 2 in [15], the set $\widehat{\mathfrak{L}}_\theta^p$ contains pseudo-diagonal channels with $d_A = 2^p$ and d_E such that $d_E^2 \geq 2^p$.

Theorem 2 Let $p > 1$ and $n > 1$ be given natural numbers, Φ_θ be an arbitrary channel in $\widehat{\mathfrak{L}}_\theta^p$ and $\delta_p = \frac{1}{2^{p-1}} \sum_{k=1}^{2^{p-1}} \left| \cot \left(\frac{(2k-1)\pi}{2^p} \right) \right| > 0$.

A) $\bar{Q}_0(\Phi_\theta^{\otimes n}) = 0$ if $|\theta| \leq \theta_n$, where θ_n is the minimal positive solution of the equation

$$2(1 - \cos(\theta/2)) + \delta_p \sin(\theta/2) = n^{-1} \ln(3/2).$$

B) If $\theta = \pm\pi/n$, then $\bar{Q}_0(\Phi_\theta^{\otimes n}) \geq p - 1$ and there exist 2^n mutually orthogonal 2^{p-1} -D error-correcting codes for the channel $\Phi_\theta^{\otimes n}$. For each binary n -tuple (x_1, \dots, x_n) , the corresponding error-correcting code is spanned by the images of the vectors

$$|\varphi_k\rangle = \frac{1}{\sqrt{2}} [|2k-1 \dots 2k-1\rangle + i |2k \dots 2k\rangle], \quad k = \overline{1, 2^{p-1}}, \quad (26)$$

under the unitary transformation $U_{x_1} \otimes \dots \otimes U_{x_n}$, where $\{|k\rangle\}$ is the canonical basis in \mathbb{C}^{2^p} , $U_0 = I_{2^p}$ and $U_1 = W_{2^p}$ (defined in (25)).

Remark 8 The constant δ_p is the Schur multiplier norm of the skew-symmetric $2^{p-1} \times 2^{p-1}$ matrix having "1" everywhere below the main diagonal. So, the sequence $\{\delta_p\}$ is nondecreasing. It is easy to see that $\delta_2 = 1$, $\delta_3 = \sqrt{2}$, $\delta_4 \approx 1.84$ and that $\delta_p = (\frac{2 \ln 2}{\pi}) p + o(p)$ for large p [10].

Note also that $\theta_n = 2 \ln(3/2) (n\delta_p)^{-1} + o(1/n)$ for large n .

Remark 9 Assertion B of Theorem 2 can be strengthened as follows:

B') If $\theta = \pm\pi/n$ then there exist 2^n mutually orthogonal projectors $P_{\bar{x}}$ of rank 2^{p-1} indexed by a binary n -tuple $\bar{x} = (x_1, \dots, x_n)$ such that

$$P_{\bar{x}} A P_{\bar{x}} = \lambda(A) P_{\bar{x}} \quad \forall A \in [\mathfrak{L}_\theta^p]^{\otimes n},$$

where $\lambda(A) \in \mathbb{C}$ does not depend on \bar{x} . $P_{\bar{x}}$ is the projector on the subspace $U_{x_1} \otimes \dots \otimes U_{x_n}(\mathcal{H}_0)$, where \mathcal{H}_0 is the linear hull of vectors (26).

This follows from the proof of Theorem 2 presented below.

Proof of Theorem 2 A) Note that $\Upsilon_{D(\theta)} - \text{Id}_{2^p}$ is the Schur multiplication by the matrix

$$-T \otimes \begin{bmatrix} u & 0 \\ 0 & u \end{bmatrix} + S \otimes \begin{bmatrix} \bar{v} & 0 \\ 0 & v \end{bmatrix},$$

where T is the $2^{p-1} \times 2^{p-1}$ matrix having "0" on the main diagonal and "1" on the other places, S is the $2^{p-1} \times 2^{p-1}$ skew-symmetric matrix having "1" everywhere below the main diagonal, $u = 1 - \Re\gamma = 1 - \cos(\theta/2)$, $v = i\Im\gamma = i \sin(\theta/2)$.

In [10], it is shown that $\|\Upsilon_S\|_{\text{cb}} = 2^{1-p}\|S\|_1 = \delta_p$. Since $\|\Upsilon_T\|_{\text{cb}} \leq 2$ and $\|\Upsilon_{A \otimes B}\|_{\text{cb}} = \|\Upsilon_A \otimes \Upsilon_B\|_{\text{cb}} = \|\Upsilon_A\|_{\text{cb}}\|\Upsilon_B\|_{\text{cb}}$, we have

$$x \doteq \|\Upsilon_{D(\theta)} - \text{Id}_{2^p}\|_{\text{cb}} \leq u\|\Upsilon_T\|_{\text{cb}} + |v|\|\Upsilon_S\|_{\text{cb}} = 2(1 - \cos(\theta/2)) + \delta_p|\sin(\theta/2)|$$

and hence $x \leq n^{-1} \ln(3/2) \leq \sqrt[p]{3/2} - 1$ if $|\theta| \leq \theta_n$.

Assume that $\bar{Q}_0(\Phi_{\theta}^{\otimes n}) > 0$ for some $\theta \in [-\theta_n, \theta_n]$. By repeating the arguments from the proof of part C of Theorem 1, we obtain

$$|\langle \psi | A | \varphi \rangle| \leq \Delta_n \quad \text{and} \quad |\langle \varphi | A | \varphi \rangle - \langle \psi | A | \psi \rangle| \leq 2\Delta_n \quad (27)$$

for some unit vectors $\varphi, \psi \in \mathbb{C}^{2^{pn}}$ and all A in the unit ball of $[\mathfrak{L}_0^p]^{\otimes n}$, where

$$\Delta_n \doteq \left\| \Upsilon_{D(\theta)}^{\otimes n} - \text{Id}_{2^{pn}} \right\|_{\text{cb}} \leq (x + 1)^n - 1 \leq 1/2.$$

Since $[\mathfrak{L}_0^p]^{\otimes n}$ is a maximal commutative $*$ -subalgebra of $\mathfrak{M}_{2^{pn}}$, Lemma 3 shows that (27) cannot be valid.

B) Let $\theta = \pm\pi/n$. To prove that the linear hull \mathcal{H}_0 of vectors (26) is an error-correcting code for the channel $\Phi_{\theta}^{\otimes n}$, it suffices, by Lemma 1, to show that

$$\langle \varphi_l | M_1 \otimes \dots \otimes M_n | \varphi_k \rangle = 0 \quad \forall M_1, \dots, M_n \in \mathfrak{L}_{\theta}^p, \quad \forall k, l$$

and that

$$\langle \varphi_l | M_1 \otimes \dots \otimes M_n | \varphi_l \rangle = \langle \varphi_k | M_1 \otimes \dots \otimes M_n | \varphi_k \rangle \quad \forall M_1, \dots, M_n \in \mathfrak{L}_{\theta}^p, \quad \forall k, l.$$

Since any matrix in \mathfrak{L}_{θ}^p can be described as $2^{p-1} \times 2^{p-1}$ matrix $[A_{ij}]$ consisting of the blocks

$$A_{ii} = \begin{bmatrix} a & b \\ b & a \end{bmatrix} \quad \forall i \quad \text{and} \quad A_{ij} = \begin{bmatrix} \bar{\gamma}_{ij} c_{ij} & d_{ij} \\ d_{ij} & \gamma_{ij} c_{ij} \end{bmatrix} \quad \forall i \neq j,$$

where $\gamma_{ij} = \exp(is_{ij}\theta/2)$, $s_{ij} = \text{sgn}(i - j)$ and a, b, c_{ij}, d_{ij} are some complex numbers, the above relations are proved by the same way as in the proof of part B of Theorem 1 (by using $\gamma_{ij}^n + \bar{\gamma}_{ij}^n = 0$).

To prove that the subspace $U_{\bar{x}}(\mathcal{H}_0)$, where $U_{\bar{x}} = U_{x_1} \otimes \dots \otimes U_{x_n}$, is an error-correcting code for the channel $\Phi_{\theta}^{\otimes n}$, it suffices to note that (25) implies that $U_{\bar{x}}^* A U_{\bar{x}} = A$ for all $A \in [\mathfrak{L}_{\theta}^p]^{\otimes n}$. \square

Corollary 3 *Let n be arbitrary and m be a natural number such that $\theta_* = \pi/m \leq \theta_n$. Then*

$$\bar{Q}_0(\Phi_{\theta_*}^{\otimes n}) = 0 \quad \text{but} \quad \bar{Q}_0(\Phi_{\theta_*}^{\otimes m}) \geq p - 1 \quad \text{and hence} \quad Q_0(\Phi_{\theta_*}) \geq (p - 1)/m.$$

There exist 2^m mutually orthogonal 2^{p-1} -D error-correcting codes for the channel $\Phi_{\theta_}^{\otimes m}$.*

Remark 10 Corollary 3 (with Proposition 2 in [15] and Remark 8) shows that for any n , there exists a channel Φ_n with $d_A = 2^p$ and arbitrary d_E satisfying the inequality $d_E^2 \geq 2^p$ such that

$$\bar{Q}_0(\Phi_n^{\otimes n}) = 0 \quad \text{and} \quad Q_0(\Phi_n) \geq \frac{p-1}{[\pi/\theta_n] + 1} = \frac{2\ln(3/2)(p-1)}{\pi n \delta_p} + o(1/n),$$

where $[x]$ is the integer part of x , and hence, we have the following lower bounds for the values $S_d(n)$ and $S_*(n)$ (introduced in 13,14)

$$S_{2^p} \geq \frac{2\ln(3/2)(p-1)}{\pi n \delta_p} + o(1/n) \quad \text{and} \quad S_*(n) \geq \frac{2\ln(3/2)(p-1)}{\pi n \delta_p}$$

(the later inequality is obtained from the former by using relation (15)).

Since $\delta_2 = 1$, the above lower bounds with $p = 2$ coincide with (16,17).

Since $\delta_3 = \sqrt{2}$, Remark 10 with $p = 3$ shows that for any n , there exists a channel Φ_n with $d_A = 8$ and $d_E = 3$ such that

$$\bar{Q}_0(\Phi_n^{\otimes n}) = 0 \quad \text{and} \quad Q_0(\Phi_n) \geq \sqrt{2} \times \frac{2\ln(3/2)}{\pi n} + o(1/n).$$

Hence,

$$S_8(n) \geq \sqrt{2} \times \frac{2\ln(3/2)}{\pi n} + o(1/n).$$

Comparing this estimation with (16), we see that the increasing input dimension d_A from 4 to 8 gives the amplification factor $\sqrt{2}$ for the quantum zero-error capacity of a channel having vanishing n -shot capacity (more precisely, for the lower bound of this capacity).

In general, Remark 10 shows that our construction with the input dimension $d_A = 2^p$ amplifies lower bound (17) for $S_*(n)$ by the factor $\Lambda_p = \delta_p^{-1}(p-1)$. By Remark 8, the nondecreasing sequence Λ_p has a finite limit:

$$\lim_{p \rightarrow +\infty} \Lambda_p = \Lambda_* \doteq \frac{\pi}{2\ln 2} \approx 2.26.$$

Hence, $\Lambda_* \approx 2.26$ is the maximal amplification factor for $S_*(n)$ which can be obtained by increasing input dimension. So, we have

$$S_*(n) \geq \Lambda_* \frac{2\ln(3/2)}{\pi n} = \frac{\log_2(3/2)}{n} \quad \forall n.$$

Unfortunately, we have not managed to show the existence of a channel with *arbitrary* quantum zero-error capacity and vanishing n -shot capacity, i.e., to prove the conjecture $S_*(n) = +\infty$ for all n . This can be explained as follows.

According to Theorem 2, if the input dimension of the channel Φ_θ increases as 2^p , then the dimension of error-correcting code for the channel $\Phi_\theta^{\otimes m}$, $\theta = \pi/m$, increases as 2^{p-1} . But simultaneously, the norm of the map $\Upsilon_{D(\theta)} - \text{Id}_{2^p}$ characterizing deformation of a maximal commutative $*$ -subalgebra increases as $\delta_p \sin(\theta/2) \sim p\theta/2$ for large p and small θ , so, to guarantee vanishing of the n -shot capacity of Φ_θ by using Lemma 3, we have to decrease the value of θ as $O(1/p)$. Since $\theta = \pi/m$, we see that $\bar{Q}_0(\Phi_\theta^{\otimes m})$ and m have the same increasing rate $O(p)$, which does not allow to obtain large values of $Q_0(\Phi_\theta)$.

Thus, the main obstacle for proving the conjecture $S_*(n) = +\infty$ consists in the unavoidable growth of the norm of the map $\Upsilon_{D(\theta)} - \text{Id}_{2^p}$ as $p \rightarrow +\infty$ (for fixed θ).

First, there was a hope to solve this problem by using a freedom in choice of the deformation map $\Upsilon_{D(\theta)}$. Indeed, instead of the matrix $D(\theta)$ introduced before the definition of \mathfrak{L}_θ^p , one can use the matrix $D(\theta, S) = [A_{ij}]$ consisting of the blocks

$$A_{ii} = \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix} \quad \forall i, \quad A_{ij} = \begin{bmatrix} \gamma & 1 \\ 1 & \bar{\gamma} \end{bmatrix} \quad \text{if } s_{ij} = -1 \quad \text{and} \quad A_{ij} = \begin{bmatrix} \bar{\gamma} & 1 \\ 1 & \gamma \end{bmatrix} \quad \text{if } s_{ij} = 1,$$

where $S = [s_{ij}]$ is any skew-symmetric $2^{p-1} \times 2^{p-1}$ matrix such that $s_{ij} = \pm 1$ for all $i \neq j$. For the corresponding subspace $\mathfrak{L}_{\theta, S}^p = \Upsilon_{D(\theta, S)}(\mathfrak{L}_0^p)$, the main assertions of Theorem 2 are valid (excepting the assertion about 2^m error-correcting codes) with the constant δ_p replaced by the norm $\|\Upsilon_S\|_{\text{cb}}$ (in our construction $S = S_*$ is the matrix having “1” everywhere below the main diagonal and $\delta_p = \|\Upsilon_{S_*}\|_{\text{cb}}$). But the further analysis (based on the results from [10]) has shown that

$$\|\Upsilon_S\|_{\text{cb}} \geq \delta_p = \|\Upsilon_{S_*}\|_{\text{cb}}$$

and hence

$$\|\Upsilon_{D(\theta, S)} - \text{Id}_{2^p}\|_{\text{cb}} \geq \|\Upsilon_{D(\theta, S_*)} - \text{Id}_{2^p}\|_{\text{cb}}$$

for any skew-symmetric $2^{p-1} \times 2^{p-1}$ matrix S such that $s_{ij} = \pm 1$ for all $i \neq j$. So, by using the above modification, we cannot increase the lower bound for $Q_0(\Phi_\theta)$. The useless of some other modifications of the map $\Upsilon_{D(\theta)}$ was also shown.

It is interesting to note that the norm growth of the map $\Upsilon_{D(\theta)} - \text{Id}_{2^p}$ is a *cost of the symmetry requirement* for the subspace \mathfrak{L}_θ^p . Indeed, if we omit this requirement, then we would use the matrix $\tilde{D}(\theta) = [\tilde{A}_{ij}]$ consisting of the blocks

$$A_{ii} = \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix} \quad \forall i \quad \text{and} \quad A_{ij} = \begin{bmatrix} \gamma & 1 \\ 1 & \bar{\gamma} \end{bmatrix} \quad \forall i \neq j,$$

for which $\|\Upsilon_{\tilde{D}(\theta)} - \text{Id}_{2^p}\|_{\text{cb}} \leq 2|\gamma - 1| \leq \theta$ for all p .

It seems that the above obstacle is technical and can be overcome (within the same construction of a channel) by finding a way to prove the equality $\bar{Q}_0(\Phi^{\otimes n}) = 0$ not using estimations of the distance between the unit balls of $[\mathfrak{L}_\theta^p]^{\otimes n}$ and of $[\mathfrak{L}_0^p]^{\otimes n}$. Anyway the question concerning the value

$$S_*(n) \doteq \sup_{\Phi} \{Q_0(\Phi) \mid \bar{Q}_0(\Phi^{\otimes n}) = 0\}$$

remains open.

Acknowledgments I am grateful to A.S.Holevo and to the participants of his seminar “Quantum probability, statistic, information” (the Steklov Mathematical Institute) for useful discussion. Special thanks to T.Shulman and P.Yaskov for the help in solving the particular questions. I am also grateful to the reviewer for the useful recommendations.

Appendix: Proof of Lemma 2

Show first that for each θ , one can construct basis $\{A_i^\theta\}_{i=1}^4$ of \mathfrak{L}_θ consisting of positive operators with $\sum_{i=1}^4 A_i^\theta = I_4$ such that:

- (1) the function $\theta \mapsto A_i^\theta$ is continuous for $i = \overline{1, 4}$;
- (2) $\{A_i^\theta\}_{i=1}^4$ consists of mutually orthogonal 1-rank projectors.

Recall that \mathfrak{L}_θ is unitary equivalent to the subspace \mathfrak{L}_θ^s defined by (9).

Denote by $\|T_\theta\|$ the operator norm of the matrix T_θ involved in (9). Note that the function $\theta \mapsto T_\theta$ is continuous, $T_0 = 0$ and $\|T_\theta\| \leq \|T_\pi\| = 2$. Let

$$\begin{aligned} \tilde{A}_1^\theta &= \begin{bmatrix} \alpha & 0 & 0 & 0 \\ 0 & \beta & 0 & 0 \\ 0 & 0 & \beta & 0 \\ 0 & 0 & 0 & \beta \end{bmatrix} - \frac{1}{4}(\alpha - \beta) T_\theta, & \tilde{A}_2^\theta &= \begin{bmatrix} \beta & 0 & 0 & 0 \\ 0 & \alpha & 0 & 0 \\ 0 & 0 & \beta & 0 \\ 0 & 0 & 0 & \beta \end{bmatrix} - \frac{1}{4}(\alpha - \beta) T_\theta, \\ \tilde{A}_3^\theta &= \begin{bmatrix} \beta & 0 & 0 & 0 \\ 0 & \beta & 0 & 0 \\ 0 & 0 & \alpha & 0 \\ 0 & 0 & 0 & \beta \end{bmatrix} + \frac{1}{4}(\alpha - \beta) T_\theta, & \tilde{A}_4^\theta &= \begin{bmatrix} \beta & 0 & 0 & 0 \\ 0 & \beta & 0 & 0 \\ 0 & 0 & \beta & 0 \\ 0 & 0 & 0 & \alpha \end{bmatrix} + \frac{1}{4}(\alpha - \beta) T_\theta \end{aligned}$$

be operators in \mathfrak{L}_θ^s , where $\beta = \min\{\frac{3}{16}, \frac{1}{4}\|T_\theta\|\}$ and $\alpha = 1 - 3\beta$. It is easy to verify that $\tilde{A}_i^\theta \geq 0$ for all i and $\sum_{i=1}^4 \tilde{A}_i^\theta = I_4$. Then $\{A_i^\theta = S\tilde{A}_i^\theta S^{-1}\}_{i=1}^4$, where S is the unitary matrix defined after (9), is a required basis of \mathfrak{L}_θ .

Let $m \geq 2$ and $\{|\psi_i\rangle\}_{i=1}^4$ be a collection of unit vectors in \mathbb{C}^m such that $\{|\psi_i\rangle\}_{i=1}^4$ is a linearly independent subset of \mathfrak{M}_m . It is easy to show (see the proof of Corollary 1 in [15]) that \mathfrak{L}_θ is a noncommutative graph of the pseudo-diagonal channel

$$\Phi_\theta(\rho) = \text{Tr}_{\mathbb{C}^m} V_\theta \rho V_\theta^*,$$

where

$$V_\theta : |\varphi\rangle \mapsto \sum_{i=1}^4 [A_i^\theta]^{1/2} |\varphi\rangle \otimes |i\rangle \otimes |\psi_i\rangle$$

is an isometry from $\mathcal{H}_A = \mathbb{C}^4$ into $\mathbb{C}^4 \otimes \mathbb{C}^4 \otimes \mathbb{C}^m$ ($\{|i\rangle\}$ is the canonical basis in \mathbb{C}^4). By property 1 of the basis $\{A_i^\theta\}_{i=1}^4$, the function $\theta \mapsto V_\theta$ is continuous.

The first part of Lemma 2 follows from this construction with $m = 2$.

To prove the second part assume that $m = 4$ and $|\psi_i\rangle = |i\rangle$, $i = \overline{1, 4}$. Property 2 of the basis $\{A_i^\theta\}_{i=1}^4$ implies

$$V_0|\varphi\rangle = \sum_{i=1}^4 \langle e_i|\varphi\rangle |e_i\rangle \otimes |i\rangle \otimes |i\rangle,$$

where $\{|e_i\rangle\}_{i=1}^4$ is an orthonormal basis in \mathbb{C}^4 . Hence, $\Phi_0(\rho) = \sum_{i=1}^4 \langle e_i|\rho|e_i\rangle \sigma_i$, $\sigma_i = |e_i \otimes i\rangle \langle e_i \otimes i|$, is a classical–quantum channel.

References

1. Holevo, A.S.: Quantum Systems, Channels, Information. A Mathematical Introduction. DeGruyter, Berlin (2012)
2. Nielsen, M.A., Chuang, I.L.: Quantum Computation and Quantum Information. Cambridge University Press, Cambridge (2000)
3. Smith, G., Yard, J.: Quantum communication with zero-capacity channels. *Science* **321**(5897), 1812–1815 (2008)
4. Cubitt, T.S., Chen, J., Harrow, A.W.: Superactivation of the asymptotic zero-error classical capacity of a quantum channel. *IEEE Trans. Inf. Theory* **57**(2), 8114–8126 (2011)
5. Cubitt, T.S., Smith, G.: An extreme form of superactivation for quantum zero-error capacities. *IEEE Trans. Inf. Theory* **58**(3), 1953–1961 (2012)
6. Duan, R.: Superactivation of zero-error capacity of noisy quantum channels. [arXiv:0906.2527](https://arxiv.org/abs/0906.2527) [quant-ph] (2009)
7. Shirokov, M.E., Shulman, T.V.: On superactivation of one-shot zero-error quantum capacity and the related property of quantum measurements. *Probl. Inform. Transm.* **50**(3), 232–246 (2014)
8. Cubitt, T.S., Elkouss, D., Matthews, W., Ozols, M., Perez-Garcia, D., Strelchuk, S.: Unbounded number of channel uses are required to see quantum capacity. [arXiv:1408.5115](https://arxiv.org/abs/1408.5115) (2014)
9. Paulsen, V.: Completely Bounded Maps and Operator Algebras. Cambridge University Press, Cambridge (2003)
10. Mathias, R.: The Hadamard operator norm of a circulant and applications. *SIAM J. Matrix Anal. Appl.* **14**(4), 1152–1167 (1993)
11. Holevo, A.S.: On complementary channels and the additivity problem. *Probab. Theory Appl.* **51**(1), 134–143 (2006)
12. Medeiros, R.A.C., de Assis, F.M.: Quantum zero-error capacity. *Int. J. Quant. Inf.* **3**(1), 135–139 (2005)
13. Duan, R., Severini, S., Winter, A.: Zero-error communication via quantum channels, non-commutative graphs and a quantum Lovasz theta function. *IEEE Trans. Inf. Theory* **59**(2), 1164–1174 (2013)
14. Knill, E., Laflamme, R.: Theory of quantum error-correcting codes. *Phys. Rev. A* **55**, 900–911 (1997)
15. Shirokov, M.E., Shulman, T.V.: On superactivation of zero-error capacities and reversibility of a quantum channel. *Commun. Math. Phys.* **335**(3), 1159–1179 (2015)
16. Cubitt, T.S., Ruskai, M.B., Smith, G.: The structure of degradable quantum channels. *J. Math. Phys.* **49**, 102104 (2008)
17. Kretschmann, D., Schlingemann, D., Werner, R.F.: The information-disturbance tradeoff and the continuity of Stinespring’s representation. [arXiv:quant-ph/0605009](https://arxiv.org/abs/quant-ph/0605009) (2006)