## **Final Exam**

1. (1 point) Is the following Boolean formula a tautology?

$$\left( (p \to (q \lor \neg r)) \land (q \to s) \land (\neg r \to (\neg s \to p)) \land (p \to \neg s) \right) \to \neg p$$

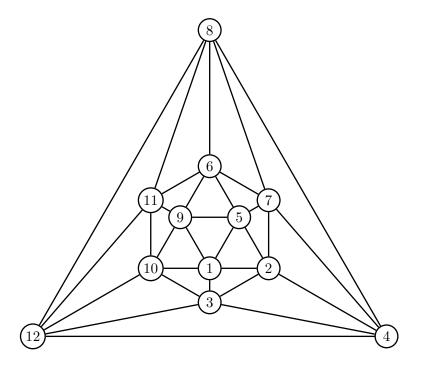
If not, provide a falsifying assignment.

2. (1 point) Is the following first-order formula satisfiable?

 $(\forall x \exists y (R(x,y) \land P(y))) \land (\forall x \exists y (R(x,y) \land \neg P(y)))$ 

If yes, provide an interpretation of predicate symbols R and P on some set (domain) M, which makes the formula true.

- 3. (1 point) Does there exist a polynomial time algorithm which, being given a graph G, yields all Euler paths in G.
- 4. (2 points) Find the minimal k for which the graph below has a correct k-coloring. (A correct k-coloring is a coloring of vertices in k colors, such that ends of each edge have different colors.) Provide a correct k-coloring and explain why (k-1) colors are not sufficient. (The coloring may be provided as a mapping, like "1 red, 2 blue, 3 green, 4 red, …")



5. (2 points) The INTPROG problem (so-called "integer programming") is formulated as follows. Given a matrix  $(a_{i,j})_{1 \le i \le n, 1 \le j \le m+1}$  of integers  $(a_{i,j} \in \mathbb{Z})$ , answer whether the system of inequations

 $\begin{cases} a_{1,1}x_1 + \dots + a_{1,m}x_m + a_{1,m+1} \ge 0\\ a_{2,1}x_1 + \dots + a_{2,m}x_m + a_{2,m+1} \ge 0\\ \dots\\ a_{n,1}x_1 + \dots + a_{n,m}x_m + a_{n,m+1} \ge 0 \end{cases}$ 

has an integer solution  $(x_1, \ldots, x_m)$  (i.e., such  $x_1, \ldots, x_m \in \mathbb{Z}$  that all inequations become true). Show that INTPROG is NP-hard by proving that 3-SAT  $\leq_m^P$  INTPROG.

6. (2 points) Suppose  $P \neq NP$ . Could there exist a polynomial-time algorithm which, being given a graph G, answers whether the number of correct 3-colorings of G is greater or equal than 3?

(One extra point is added just for the fact of participating in the exam. Participating means submitting something within the scope of the deadline.)